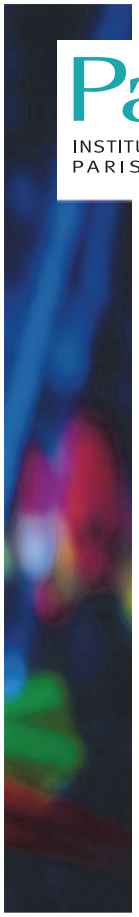
Visual-based event mining in social media

Riadh Trad



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Découverte d'événements par contenu visuel dans les médias sociaux

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T H È S E

ABSTRACT

Visual-based event mining in Social Media

Trad Mohamed Riadh

Broadly understood, events are things that happen, things such births, weddings and deaths. Among these moments, are those we record and share with friends. Social media sites, such as Flickr or Facebook, provide a platform for people to promote social events and organize content in an event centric manner.

The ease of capturing and publishing on social media sites however, has led to significant impact on the overall information available to the user. The number of social media documents for each event is potentially very large and is often disseminated between users. Defining new methods for organizing, searching and browsing media according to real-life events is therefore of prime interest for maintaining a high-quality user experience.

While earlier studies were based solely on text analysis and essentially focused on news documents, more recent work has been able to take advantage of richer multimedia content available, while having to cope with the challenges that such a benefit entails. And as the amount of content grows, research will have to identify robust ways to process, organize and filter that content. In this dissertation we aim to provide scalable, cloud oriented techniques for organizing social media documents associated with events, notably in scalable and distributed environments.

To identify event content, we develop a visual-based method for retrieving events in photo collections, typically in the context of User Generated Content.

Given a query event record, represented by a set of photos, we aim at retrieving other records of the same event, typically generated by distinct users.

Matching event records however, requires defining a similarity function that captures the multi-facet similarity between event records. Although records of the same event often exhibit similar facets, they may differ in several aspects. Records of the same event, for instance, are not necessarily located at the same place (*e.g.* an eclipse, tsunami) and can be recorded at different times (*e.g.* during a festival). Importantly, we show how using visual content as complementary information might overcome several limitations of state-of-the-art approaches that rely only on metadata.

The number of social media documents for each event is potentially large. While some of their content might be interesting and useful, a considerable amount might be of little value to people interested in learning about the event itself. To avoid overwhelming users with unmanageable volumes of event information, we present a new collaborative content-based filtering technique for selecting relevant documents for a given event. Specifically, we leverage the social context provided by the social media to objectively detect moments of interest in social events. Should a sufficient number of users take a large number of shots at a particular moment, then we might consider this to be an objective evaluation of interest at that moment.

With our record-based event retrieving paradigm, we provide novel approaches for searching, browsing and organizing events in social media. Overall, the work presented in this dissertation provides an essential methodology for organizing social media documents that reflect event information, with a view to improving browsing and search for social media event data.

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I am also grateful to the members of my oral defense committee : Matthieu CORD, M. Bernard MERIALDO, Denis TEYSSOU, Frédéric PRECIOSO and Nicu SEBE, for their time and insightful questions.

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TRAD Mohamed Riadh

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March on. Do not tarry. To go forward is to move toward perfection. March on, and fear not the thorns, or the sharp stones on life's path.

Gibran Khalil Gibran

Résumé

1 Motivations

L'évolution du web, de ce qui était typiquement connu comme un moyen de communication à sens unique en mode conversationnel, a radicalement changé notre manière à traiter l'information.

Des sites de médias sociaux tels que Flickr et Facebook, offrent des espaces d'échange et de diffusion de l'information. Une information de plus en plus riche, abondante mais aussi personnelle, *i.e.* capturée par l'utilisateur, et qui s'organise, le plus souvent, autour d'événements de la vie rélle.

Ainsi, un événement peut être perçu comme un ensemble de vues personnelles et locales, capturées par les utilisateurs ayant pris part à l'événement. Identifier ces différrentes instances permettrait, dès lors, de reconstituer une vue globale de l'événement. Plus particulièrement, lier différentes instances d'un même événement profiterait à bon nombre d'applications tel que la recherche, la navigation ou encore le filtrage et la suggestion de contenus.

2 Problématiques

L'objectif principal de cette thèse est l'identification du contenu multimédia, associé à un événement dans de grandes collections d'images.

Plus particulièrement, on s'intéresse au contenu généré par l'utilisateur et publié dans les médias sociaux. Un tel contenu est le plus souvent, diffus, partagé par différents utilisateurs, il peut être hétérogène, bruité et pour la plus part non annoté.

Pour mieux illustrer la motivation derrière l'identification du contenu d'événements dans les médias sociaux, considérons une personne souhaitant assister au festival de musique "Rock en Seine" dans le parc du chaâteau de Saint-Cloud. Dans ce sens, et avant de prendre sa décion, elle entreprend une recherche concernant les éditions des années précédentes. Le site de l'événement contient des informations de base, le programme, la billetterie et même si il y figure des médias des éditions précédentes, elle ne reflètent pas l'ambiance du festival. Cette couverture large rend les sites de médias sociaux une source inestimable d'informations.

Après avoir assisté à l'événement, un utilisateur peut vouloir prolonger son expérrance en visionnant des médias capturés par d'autres utilisateurs. En téléchargeant le contenu qu'il a capturé lors de l'événement un utilisateur peut avoir accès au contenu généré par d'autre participants. L'utilisateur peut alors revivre l'événement en naviguant dans les photos prises par d'autres utilisateurs, il peut par ailleurs enrichir sa propre collections de médias ou encore contribuer à enrichir le contenu global disponible sur le web.

Dans un contexte plus professionnel, détecter automatique le fait qu'un grand nombre d'utilisateurs s'intéressent à un même événement peut être utile pour orienter les journalistes vers des événements imprévus, ou encore approcher des utilisateurs pour récupérer du contenu.

Dans de tels scénarios, l'information spatiale et temporelle associée au contenu a un rôle majeur. Cependant, dans ce qui suit, nous montrons que l'utilisation du contenus visuel est indispensable. En effet, des instances distinctes d'un événement ne sont pas nécessairement localisées au même endroit et peuvent être enregistrées à des moments différents. Certains événement, naturels par example, on des étendues spatiales vastes, dans ces cas, l'utilisations des métadonnées n'est pas assez discriminante. Ceci est encore valable pour les événement colocalisés, typiquement dans des lieux très fréquentés comme les gares, les centres commerciaux ou les lieux touristiques. Dans de tels environnements, les médias générés portent les même signatures temporelles et spatiales, nonobstant le fait qu'ils soient associés à des événements distincts. Plus généralement, plusieurs instances d'un événement peuvent être enregistrées à des moments différents. Enfin, les informations spatiales et temporelles ne sont pas toujours disponibles ou pourraient être biaisées.

3 Contributions

3.1 Événement et instances d'événement

Dans le premier chapitre de cette thèse, nous dressons un état de l'art sur la définition de la notion d'événement dans la littérature. Bien que la plus part de ces définitions s'accordent sur les principales facettes pour la caractérisation d'événements (intervenants, lieux, temps, ...), elles ignorent le contexte social associé aux médias décrivant un événement.

Ainsi, et partant de l'hypothèse cadre de cette thìse que chaque média décrivant un événements est généré par un utilisateur, un événement peut être vu comme l'ensemble de ces vues personnelles et locales, capturées par les différents utilisateurs.

Dans ce qui suit, on désigne par enregistrement l'ensemble des médias générés

par un utilisateur au cours d'un événement.

3.2 Recherche d'événements par similarité visuelle

Afin d'identifier les différentes instances d'un même événement, nous proposons une méthode de recherche par similarité visuelle dans des collections d'images.

Étant donné un enregistrement requête d'images d'un événement, notre méthode vise à identifier d'autres enregistrements du même événement, typiquement générés par d'autres utilisateurs. L'appariement d'instances d'événements requiert néansmoins la définition d'une mesure de simimilarité capable de capturer la similarité multi-facette entre les enregistrements.

Cependant, bien que différents enregistrement d'un même événement présentent le plus souvent des caractéristiques similaires, ils peuvent néansmoins être différents sur certains aspects. Les enregistrements d'un même événement, par exemple, ne sont pas nécessairement localisés au même endroit (e.g. une éclipse, un tsunami) et peuvent être capturés à des moments différents (e.g. lors d'un festival). Par ailleurs, les informations tel que les coordonnées spatiales ou temporelles sont, le plus souvent manquante ou biaisées. Ceci limite leur utilisation.

Afin de palier à de telles limites, nous proposons une stratégie en deux étapes combinant à la fois continu visuel et le contexte associé aux médias jusqu'alors non exploités. Une première étape vise à identifier un premier ensemble d'enregistrements, visuellement similaires à l'enregistrement constituant requête. Une deuxième étape vise à filtrer et à reclasser les enregistrements via un recordage spatio-temporel avec l'enregistrement requète.

Identifier les différentes instances d'un événement peut s'avérrer utile pour diverses applications notamment pour l'identification de contenu d'un événement dans une collection d'images ou encore la génération automatique de contenu et

nécessitent le plus souvent la construction de graphes de similarités entre les différents enregistrements.

3.3 Construction scalable et distribuée du graphe de similarité visuelle

L'appariement d'enregistrements d'événements requiert l'appariement par contenu visuel entre images appartenants à différents enregistrements.

Une solution naïve, serait de chercher les k-images les plus similaires visuellement à chaque image de l'enregistrement requête. Une telle approche peut s'avérer coûteuse si l'on considérait la construction du graphe de similarité sur l'ensemble des enregistrements d'une collection d'images.

Le principal problème de la construction d'un tel graphe est le temps de calcul. La complexité de l'approche naïve est certes linéaire en nombre d'images, mais la recherche reste coûteuse, à moins de fortement dégrader la qualité au profit de la vitesse en effectuant des recherches approximatives. Là encore, le coût de la recherche reste tributaire du choix des fonctions de hachages. Par ailleurs, de telles approches restent difficilement distribuables du fait qu'elle requièrent la duplication des données sur les unités de traitement et le plus souvent leur chargement en mémoire et donc, passent difficilement à l'échelle.

D'autres approchent sont alors envisageables. Dans [16] , Chen et al. proposent de subdiviser l'ensemble des données puis de construire les graphes associés pour enfin les combiner en une solution finale au problème. Ici, le problème réside dans le choix des différentes partitions.

Dans [31], Dong et al. proposent de partir d'une solution aléatoire et partant du principe qu'un plus proche voisin d'un plus proche voisin est potentiellement un plus proche voisin, converger vers une solution au problème en un nombre faible

d'itérations. Cependant, un tel algorithme reste difficilement distribuable.

La solution que nous proposons découle d'une analyse ascendante du problème. Une solution à la fois distribuable et scalable exige de petites unités de traitement.

Par ailleurs, calculer la similarité entre les objets peut s'avérer coûteux. Ceci est d'autant plus vrai lorsqu'il s'agit de traiter de très grandes collections d'objets. Ici nous considérons que le nombre de fois que deux objets sont mappés dans une même bucket est une estimation pertinente de la similarité entre ces deux objets.

3.4 Sélection de contenu

Le nombre de documents associés à un événement dans les médias sociaux est potentiellement grand. Filtrer un tel contenu peut s'avérer bénéfique pour des applications tel que la recherche, la navigation ou encore l'organisation de contenus.

Plus particulièrement, nous nous intéressons à la sélection de contenus pertinents pour la génération automatique de résumés d'événements.

Plusieurs travaux se sont intéressés à définir des mesures capables de capturer l'importance d'un document de manière objective. Plustôt que de s'attarder à la définition et à l'évaluation d'une telle mesure, nous considérons que le nombre de photos capturées, se rapportant à une même scènes, par différents utilisateurs comme une mesure objective de son importance.

Une approche naïve consiste à compter le nombre d'images prises sur un intervalle de temps donné et localisées avec des coordonnées spatiales bien déterminése. Cependant, les métadonnées associées au contenu sont souvent absentes ou biaisées. On pourrait par ailleurs compter le nombre d'images visuellement similaires. Cependant le contenu visuel est le plus souvent non discriminant. Pour pallier à cette limite, nous effectuons un recordage spatio-temporel des enregistrements

d'un même événement, puis comptons le nombre d'image visuellement similaires entre les différents enregistrements (i.e. qui contribuent à l'appariement de deux enregistrements).

Génération automatique de résumé

Nous ramenons le problème de génération automatique de résumé à celui de produire un classement sur les documents d'un événement. L'ensemble des images sélectionnées est ensuite traité afin d'en éliminer les doublons.

Alternativement, l'ensemble des images est filtré pour produire des résumés personnalisés en fonction de la qualité des images ou encore les droits associés.

Suggestion de contenu

Un contenu est dit intéressant, d'un point de vue utilisateur, s'il renseigne sur des aspects de l'événement, autres que ceux capturés par l'utilisateur. Suggérer du contenu, reviendrait, dès lors, à proposer du contenu pertinent, visuellement différent du contenu capturé par l'utilisateur.

Organisation de la thèse

Cette thèse s'organise comme suit :

Dans le **Chapitre 2**, nous examinons différentes définitions de la notion d'événement dans la littérature, puis proposerons une définition alternative qui tient compte à la fois du contenu visuel et du contexte.

Le **Chapitre 3** présente notre méthode de recherche d'événements basée sur le contenu visuel dans les collections d'images.

Le **Chapitre 4** présente notre approche pour la construction scalable et distribuée

des Graphe des K plus proches voisins et son implémentation dans le framework Hadoop.

Dans le **Chapitre 5**, nous présentons notre méthode collaborative pour la sélection de contenu pertinent dans de grandes collections d'images. Plus particuliérement, nous nous intŕesserons aux problèmes de génŕation automatique de résumés d'événements et suggestion de contenus dans les médias sociaux.

Le **Chapitre 6** dresse un état-de-l'art des problématiques abordées dans cette thèse.

Le bilan des contributions, la conclusion et les perspectives sont présentés dans le **Chapitre 7**.

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Chapitre 1

General Introduction

Problem Statement

Social Media sites such as Flickr and Facebook, have changed the way we share and manage information within our social networks. The shift on the Web, from what was typically a one-way communication, to a conversation style interaction has led to many exciting new possibilities.

The ease of publishing content on social media sites brings to the Web an ever increasing amount of user generated content captured during, and associated with, real life events. Social media documents shared by users often reflect their personal experience of the event. Hence, an event can be seen as a set of personal and local views, recorded by different users. These event records are likely to exhibit similar facets of the event but also specific aspects. By linking different records of the same event occurrence we can enable rich search and browsing of social media events content. Specifically, linking all the occurrences of the same event would provide a general overview of the event. In this dissertation we present a content-based approach for leveraging the wealth of social media documents available on the Web for event identification and characterization.

To better illustrate the motivation behind event content identification in social media, consider a person who is planning to attend the "Rock en Seine" annual music Festival in Château de Saint-Cloud's Park. Before buying a ticket, the person could do some research upon which he will make an informed decision. The event's website contains basic information about the festival and the tickets available. Although the event website contains stage pictures and videos of prior instances of the event, they do not reflect the general atmosphere of the event. User-generated content may, however, provide a better overview of prior occurrences of the event from an attendee's perspective. Such wide coverage makes social media sites an invaluable source of event information.

After attending the event, the user may be interested in retrieving additional media associated to the event. By simply uploading his/her own set of event pictures a user might for example access to the community of the other event's participants. The user can then revive the event by browsing or collecting new data complementary to his/her own view of the event. If some previous event's pictures were already uploaded and annotated, the system might also automatically annotate the set or suggest some relevant tags to the user.

In a more professional context, automatically detecting the fact that a large number of amateur users did record data about the same event would be very helpful for professional journalists in order to cover breaking news. Finally, tracking events across different media also has a big potential for historians, sociologists, politicians, etc.

Of course, in such scenarios, time and geographic information provided with the contents has a major role to play. Our claim is that using visual content as complementary information might solve several limitations of approaches that rely only on metadata. First of all, distinct instances of the same event are not necessarily located in the same place or can be recorded at different times. Some

events might, for example, have wide spatial and temporal extent such as a volcano eruption or an eclipse, so that geo-coordinates and time stamps might be not discriminant enough. This lack of discrimination can be problematic even for precisely located events, typically in crowded environments such as train stations, malls or tourist locations. In such environments, many records might be produced at the same time and place while being related to very distinct real-world events. Furthermore, in a broader interpretation of the event concept, several instances of an event might be recorded at different times. Finally, location and time information is not always available or might be noisy. The Flickr data used in our experiments notably does not contain any geographic information and contains noisy time information.

Our problem is more similar to the MediaEval Social Event Detection Task [1], which aims to develop techniques to discover events and detect media items that are related to either a specific social event or an event-class of interest. However, our approach exhibits some fundamental differences from the traditional social event detection task that originate from the focus on content distribution across event participants.

To match event occurrences in social media, we develop a new visual-based method for retrieving events in photo collections, typically in the context of User Generated Content. Given a query event record, represented by a set of photos, our method aims at retrieving other records of the same event, typically generated by distinct users. Similarly to what is done in state-of-the-art object retrieval systems, we propose a two-stage strategy combining an efficient visual indexing model with a spatiotemporal verification re-ranking stage to improve query performance. Visual content is used in a first stage to detect potential matches, while geo-temporal metadata are used in a second stage to re-rank the results and therefore estimate the spatio-temporal offset between records.

1. http ://www.multimediaeval.org/mediaeval2011/SED2011/

The number of social media documents for each event is potentially very large. While some of their content might be interesting and useful, a considerable amount might be of little value to people interested in learning about the event itself. To avoid overwhelming users with unmanageable volumes of event information, we present a new collaborative content-based filtering technique for selecting relevant documents for a given event. Specifically, we leverage the social context provided by the social media to objectively detect moments of interest in social events. Should a sufficient number of users take a large number of shots at a particular moment, then we might consider this to be an objective evaluation of interest at that moment.

As the amount of user generated content increases, research will have to develop robust ways to process, organize and filter that content. In this dissertation we present scalable techniques for organizing social media documents associated with events, notably in distributed environments

Contributions

The research described in this thesis led to the following results :

1. A new visual-based method for retrieving events in photo collections.
2. A scalable and distributed framework for Nearest Neighbors Graph construction for high dimensional data.
3. A collaborative content-based filtering technique for selecting relevant social media documents for a given event.

Outline

This chapter informally introduces the questions investigated in this thesis. The remaining part of this thesis is structured as follows :

Chapter 2 discusses several alternative definitions of events in the literature and provides the event definitions that we use in this dissertation.

Chapter 3 presents our new visual-based method for retrieving events in photo collections.
Chapter 4 describes our large scale K-Nearest Neighbors Graph construction technique that we considered for event graph construction.

Chapter 5 presents our collaborative content-based content selection technique. Specifically, we address the problem of event summarization and content suggestion in social media.

Chapter 6 reviews the literature that is relevant to this dissertation.

Chapter 7 presents our conclusions and discusses directions for future work.

Publications

The work presented in this manuscript has led to the following publications :

Conferences

— M. R. Trad, A. Joly, and N. Boujemaa. Large scale visual-based event matching. In *Proceedings of the 1st ACM International Conference on Multimedia Retrieval*, ICMR '11, pages 53 :1–53 :7, New York, NY, USA, 2011. ACM.

— M. R. Trad, A. Joly, and N. Boujemaa. Distributed knn-graph approximation via hashing. In *Proceedings of the 2nd ACM International Conference on Multimedia Retrieval*, ICMR '12, pages 43 :1–43 :8, New York, NY, USA,

2012. ACM.

— R. Trad, Mohamed, A. Joly, and N. Boujemaa. Distributed approximate KNN Graph construction for high dimensional Data. In *BDA - 28e journées Bases de Données Avancées - 2012*, Clermont-Ferrand, France, Oct. 2012.

— R. Trad, Mohamed, A. Joly, and B. Nozha. Large scale knn-graph approximation. In *The 3rd International Workshop on Knowledge Discovery Using Cloud and Distributed Computing Platforms (KDCloud, 2012) jointly held with IEEE ICDM 2012*, Brussels, Belgium, December 2012.

Chapitre 2

Events in Social Media

Introduction

Broadly understood, events are things that happen, things such births and deaths, celebrations and funerals, elections and impeachments, smiles, shows and explosions. Yet although the definition and characterization of an "event" has received substantial attention across various academic fields [105, 13], it is not clear what precisely constitutes an event.

Often, an event is described as an abstract concept [13], or defined within the context of a very specific domain. In this chapter, we survey a number of definitions from various domains, particularly that of social media (Section 1) and draw on them to define an event with respect to our work (Section 2).

1 Events in the literature

While previous research on events has focused solely on textual news documents [61, 63], more recent efforts have been concerned with a richer content [66, 82, 90]. In this section, we look at various efforts to define events in the context of social media through four different tasks : Topic Detection and Tracking in news documents (Section 1.1), event extraction from unstructured text

(Section 1.2), multimedia event detection (Section 1.3) and social event detection (Section 1.4).

1.1 Topic Detection and Tracking

The Topic Detection and Tracking (TDT) initiative was first intended to explore techniques for identifying events and tracking their reappearance and evolution in a text document stream. Within the TDT context, an event was initially defined as "some unique thing that happens at some point in time" [3]. This definition was further extended to include location as well [104], defining an event as "something that happens at some specific time and place".

Under this definition, the World Trade Center attacks that took place on September 11, 2001 is an event. However, the media also reported the subsequent collapse of the World Trade Center towers. Here, it is unclear whether the events should be considered as separate events or whether they form part of one single event.

To address such an ambiguity, an amended definition was proposed in [2] stating that an event is "a specific thing that happens at a specific time and place along with all necessary preconditions and unavoidable consequences". Although this definition makes some clarifications regarding event boundaries, it does not cover all possible types of event, since some of the necessary preconditions and unavoidable consequences may be ambiguous, unknown or subject to debate.

Although the TDT-inspired definitions of an event introduce some useful concepts, they do not cover all possible types of events.

1.2 Event Extraction

Event extraction from unstructured data such as news messages is a task that aims at identifying instances of specific types of events, and their associated attributes [42].

The Automatic Content Extraction [1] (ACE), for instance, defines an event "as a specific occurrence involving participants". However, rather than defining all possible events abstractly, events are defined according to their expression in unstructured text and provides a set of corresponding predefined templates along with their specific pre-defined attributes (time, place, participants, etc.). An event is identified via a keyword trigger (i.e. the main word which most clearly expresses an event's occurrence) and detects the corresponding set of attributes. A template of the "attack" event subtype applied to the sentence "A car bomb exploded Thursday in a crowded outdoor market in the heart of Jerusalem, killing at least two people, police said." is presented in Table 2.1.

Attribute	Description	Example
Attacker	The attacking/instigating agent	demonstrators
Target	The target of the attack	Israeli soldiers
Instrument	The instrument used in the attack	stones and empty bottles
Time	When the attack takes place	yesterday
Place	Where the attack takes place	a Jewish holy site at the town's entrance

TABLE 2.1 – Attack event template and sample extracted attributes [2].

The ACE event definition makes the implicit assumption that events should have one or more participants. Yet, not all events have a clearly defined set of participants, thus limiting its practical use. The same remark applies for the time and place attributes. Although they were not mentioned in this definition, they are also present in almost al of the predefined templates.

As opposed to the TDT inspired definitions, the ACE-inspired definition is specific and restricted to a small class of events. Besides, this definition is only applicable to supervised event detection tasks, where the classes of events are known a priori. One drawback is that events such as Festivals and Concerts cannot be represented since there are no corresponding templates.

1. http ://projects.ldc.upenn.edu/ace/docs/English-Events-Guidelines_v5.4.3.pdf
2. http ://projects.ldc.upenn.edu/ace/docs/English-Events-Guidelines_v5.4.3.pdf

1.3 Multimedia Event Detection

Multimedia Event Detection (MED) as part of the TREC Video Retrieval Evaluation [3] aims to develop event detection techniques to enable a quick and accurate search for user-defined events in multimedia collections. An event, according to the MED [4] 2010, is "an activity-centered happening that involves people engaged in process-driven actions with other people and/or objects at a specific place and time".

Contrarily to the above described event detection tasks, the use of media associated human-annotated textual context features (e.g., title, tags) is not allowed. Each event has a corresponding Òevent kitÓ consisting of a name, a definition, an evidential description (Table 2.2) and a set of illustrative video examples.

Event Name	Assembling a shelter
Definition	One or more people construct a temporary or semi-permanent shelter for humans that could provide protection from the elements
Evidential Description	primarily outdoor settings during the day or night
Scene	
Objects/People	cutting and digging tools, tent poles and flys, tents, stakes, tree limbs, tree branches
Activities	clearing land, cutting trees and branches, gathering flooring material, assembling a tent, lashing limbs together, staking down poles
Exemplars	

TABLE 2.2 – Example of an "event kit" for the MED task [5].

1.4 Social Event Detection

Similar to the MED event detection task, the Social Event Detection (SED) task aims to discover events and their related media items. Extracting such events from multimedia content has been the focus of numerous efforts as part of the

3. http ://trecvid.nist.gov/
4. http ://www.nist.gov/itl/iad/mig/med10.cfm
5. http ://projects.ldc.upenn.edu/havic/MED10/EventKits.html

MediaEval 2011 Social Event Detection (SED) [75] task. The SED guidelines [6] define the social aspect of an event but do not provide a precise definition of the event. According to SED, social events are "events planned by people and attended by people". It also requires the social media be "captured by people".

Although the MediaEval 2011 Social Event Detection task did not provide a precise event definition, the proposed methods only exploited some known event features, namely, event title, venue and time. These attributes were also used in [102, 91] to define an event according to its context, a set of facets (image, who, when, where, what) that support the understanding of everyday events.

2 Events in social media

Going back to the September 11 example, according to some definitions, it might be considered to be an event, but it is not an event in social media until it has a corresponding realization in social media documents. Instead of providing an abstract, ambiguous, or arguable definition of an event, we extend previous definitions to include at least one single document. In our work, we focus solely on user generated pictures of events. Formally, we define an event as :

Definition 1 *An event in social media is a real world occurrence e with (1) an associated time period T_e, (2) a set of documents D_e about the occurrence, and (3) one or more features that describe the occurrence.*

The time period T_e in our definition delimits the event occurrence in time. Several records of the same event might however be time coded differently (i.e. time shifts, wide and temporal extent of the event), and so time offsets should therefore be tolerated. Moreover, documents related to an event could be produced before or after its occurrence. For instance, in our "Rock in Rio Festival 2012" example, a photograph of a participant at the Lisbon Portela Airport represents the author's experience in the context of the event and will, therefore, be associated with the

6. http ://www.multimediaeval.org/mediaeval2011/SED2011/

event for our purposes. Here, it is worth noticing that upload time often differs from the event time period and may not preserve the temporal coherence of the documents.

The *document set* in our definition (Definition 1) refers to a set of social media documents, which can be represented using a variety of associated context *features* (e.g., title, tags, signatures, timestamp). Within the context of social media, each document is typically associated to at least one user, the one who first uploaded the picture. A single image, however, may tell different stories, with different people through different events. Hence, we associate each image with the user who shared the document, regardless of its provenance.

The *features set* in our definition may include information such as the event title, location or the set of participants. As discussed above, such a definition is prone to ambiguity as it does not include all possible types of events. However, we believe that such attributes can be relaxed when considering visual information about the event. Thus, according to our event definition, events in social media include widely known occurrences such as earthquakes, and also local and private events such as festivals and weddings.

Most often, images shared by users reflect their personal experience of the event. In this connection, an event occurrence can be seen as a set of personal and local views, recorded by different users. These event records are likely to exhibit similar facets of the event but also specific aspects. Linking different records of the same occurrence would provide a general overview of the event.

Definition 2 *An event record is a set of images (1) shared by a user, (2) reflecting his/her own experience of a given event.*

Non-event content, of course, is prominent on social media sites. In our work, we make the assumption that event-related documents are shared in separate albums (i.e. records). However, our approach can generally be extended to handle less structured content. In [82] for instance, the authors present an approach for grouping photos that belong to the same event within Facebook albums using

clustering algorithms on their upload time.

3 Related tasks

Considering the fact that event related documents are often distributed among different users (i.e. event records), we extend existing tasks to support local experiences of the event.

3.1 Event matching

Given a query event record, represented by a set of photos, the event matching task aims to retrieve other records of the same event, typically generated by distinct users. Linking different occurrences of the same event would enable a number of applications such as search, browsing and event identification.

Matching event records, requires defining a similarity function that measures the multi-facet similarity between event records. Although records of the same event often exhibit similar facets, they may differ in several aspects. Records of the same event, for instance, are not necessarily located at the same place (eclipse, tsunami) and can be recorded at different times (festival). This lack of discrimination can be problematic even for precisely located events, typically in crowded environments such as train stations, shopping malls or touristic areas. In such environments, many records might be produced at the same time and place while being related to very distinct real-world events. Designing such a similarity function is, thus, a tricky task.

In Chapter 3, we show how using visual content as complementary information might solve several limitations of approaches that rely only on metadata. To the best of our knowledge, none of the existing studies have addressed the problem of linking different occurrences of the same real-word event. This is in contrast to the literature which considers an event as a set of documents, regardless of their social context. The state-of-the-art presented in Section 1 is related to the more general problem of identifying documents of the same event, i.e. the different occurrences

of the event.

According to our definition, event-related records can be seen as a connected subgraph of the records Nearest Neighbors Graph, ideally a complete graph of the event records. In Chapter 4, we present a distributed framework for approximate K-NNG construction to address the more general problem of identifying documents of the same event in very large scale contexts.

3.2 Content Selection

Events in social media are likely to have huge amounts of associated content. For instance, as of October 2012, the 2012 Rock in Rio Festival has over 6,000 associated Flickr photos. This is not limited to world renowned events, but is also true for smaller events that could feature up to dozens to hundreds of different documents. Being able to rank and filter event content is crucial for a variety of applications such as content suggestion and event summarization.

In this connection, the content selection task aims at selecting relevant documents for people who are seeking information about a given event. Nevertheless, selecting the most interesting images often involves some decision-making, based on various criteria.

Most state-of-the-art approaches reduce the problem of selecting images from photo collections to an optimization problem under quality constraints. Choosing the right combination of these criteria is a challenging problem in itself.

Most significantly, with a few exceptions, existing work often ignores the social context of the images. Obviously, should a sufficient number of users take a large number of shots at a particular moment, then we might consider this to be an objective evaluation of interest at that moment.

In Chapter 5, we present a visual-context based approach for detecting moments of interest and subsequently, interesting shots (Section 1.1). We then address the problem of content suggestion and event summarization separately.

Content Suggestion

The content suggestion task is related to the content selection task, but instead of selecting a set of potentially interesting documents, it aims to present a given user only documents that provide additional information about the event.

Recently, there has been a body of work on content suggestion (Section 2) but none has considered the use of the social context provided by the media sites. Here, we link the content suggestion problem to the previously introduced event matching task to present a novel approach for suggesting and sharing complementary information between people who attended or took part in the same event (Section 1.3).

Event Summarization

The event summarization task aims to construct a minimal yet global summary of the event.

The problem of summarizing event-related documents has been extensively addressed across different domains (Section 2), from free text documents (system logs) to more richer data representations (images, sound and videos). Many complex systems, for instance, employ sophisticated record-keeping mechanisms that log all kinds of events that occurred in the systems.

Still, event related documents in social media are often produced and even uploaded by distinct users resulting in data redundancy (London 2012 Olympic Opening Ceremony shots shared by different people) and duplication (the same picture shared by distinct users). In Section 1.2, we show how to leverage document redundancy between distinct users to produce high quality event summaries.

4 Conclusion

Although information such as location and time eliminate ambiguity in event definitions, they are also restrictive as they do not apply to all possible types of

events. Our claim is that using visual content as complementary information might relax some conditions on such attributes. This is particularly true in social media, where textual data are very rare, and metadata noisy but where visual content is abundant.

Chapitre 3

Event Identification in Social Media

Events are a natural way for referring to any observable occurrence grouping people in a specific time and place. Events are also observable experiences that are often documented by people through different media. This notion is potentially useful for connecting individual facts and discovering complex relationships. Defining new methods for organizing, searching and browsing media according to real-life events is therefore of prime importance for ultimately improving the user experience.

In this chapter we introduce a new visual-based method for retrieving events in photo collections, typically in the context of User Generated Contents. Given a query event record, represented by a set of photos, our method aims to retrieve other records of the same event, typically generated by distinct users. In Section 1, we first discuss the interest and implications of such a retrieval paradigm. Section 2 introduces our new visual-based event matching technique and its implementation in the MapReduce framework (Section 3). Section 4 reports results on a large dataset for distinct scenarios, including event retrieval, automatic annotation and tags suggestion. The bulk of this chapter appeared in [94].

1 Towards event centric content organization in social media

Multimedia documents in User Generated Content (UGC) websites, as well as in personal collections, are often organized into events. Users are usually more likely to upload or gather pictures related to the same event, such as a given holiday trip, a music concert, a wedding, etc. This also applies to professional contents such as journalism or historical data that are even more systematically organized according to hierarchies of events.

Given a query event record represented by a set of photos, our method aims to retrieve other records of the same event, notably those generated by other actors or witnesses of the same real-world event. An illustration of two matching event records is presented in Figure 3.1. It shows how a small subset of visually similar and temporally coherent pictures might be used to match the two records, even if they include other distinct pictures covering different aspects of the event. Application scenarios related to such a retrieval paradigm are numerous. By simply uploading their own record of an event users might, for example, gain access to the community of other participants. They can then *revive* the event by browsing or collecting new data complementary to their own view of the event. If some previous event's records had already been uploaded and annotated, the system might also automatically annotate a new record or suggest some relevant tags. The proposed method might also have nice applications in the context of citizen journalism. Automatically detecting the fact that a large number of amateur users did indeed record data about the same event would be very helpful for professional journalists in order to cover breaking news. Finally, tracking events across different media has a big potential for historians, sociologists, politicians, etc.

Of course, in such scenarios, time and geographic information provided with the contents has a major role to play. Our claim is that using visual content as complementary information might overcome several limitations of approaches that

FIGURE 3.1 – Two events records of an Alanis Morissette concert

rely only on metadata. First of all, distinct records of the same event are not necessarily located at the same place or can be recorded at different times. Some events might, for example, have wide spatial and temporal coverage such as a volcano eruption or an eclipse, so that geo-coordinates and time stamps might not be sufficiently discriminant. This lack of discrimination can be problematic even for precisely located events, typically in crowded environments such as train stations, malls or tourist locations. In such environments, many records might be produced at the same time and place while being related to very distinct real-world events. Furthermore, in a wider meaning of the *event* concept, several instances of an event might be recorded at different times, e.g. periodical events or events such as "a trip to Egypt" illustrated in Figure 3.2. Finally, location and time information is not always available or might be noisy. The Flickr dataset used in the experiments reported in this chapter notably does not contain any geographic information and contains noisy time information (as discussed in Section 4).

Finally, our work is, to some extent, related to object retrieval in picture collections. Our method is indeed very similar to state-of-the-art large-scale object retrieval methods combining efficient bag-of-words or indexing models with a spatial verification re-ranking stage to improve query performance [79, 53]. We might give the following analogy : images are replaced by event records (picture sets), local visual features are replaced by global visual features describing each picture of a record, spatial positions of the local features are replaced by the geo-coordinates

FIGURE 3.2 – Two records of the event "a trip in Egypt"

and time stamps of the pictures. Matching spatially and temporally coherent event records is finally equivalent to retrieving geometrically consistent visual objects.

2 Visual based Event Matching

We first describe the proposed method in the general context of event records composed of a set of geo-tagged and time coded pictures. We further restrict ourselves to time coded only pictures since our experimental dataset did not include geo-tags.

We consider a set of N event records E_i, each record being composed of N_i pictures I_j^i captured from the same real-world event. Each picture is associated with a geo-coordinate $\mathbf{x}_j^i$ and a time stamp t_j^i resulting in a final geo-temporal coordinate vector $\mathbf{P}_j^i = (\mathbf{x}_j^i, t_j^i)$. The visual content of each image I_j^i is described by a visual feature vector $\mathbf{F}_j^i \in \mathbb{R}^d$ associated with a metric $d : \mathbb{R}^d \times \mathbb{R}^d \to \mathbb{R}$.
Now let E_q be a query event record represented by N_q pictures, with associated

visual features $\mathbf{F}_j^q$ and geo-temporal metadata $\mathbf{P}_j^q$. Our retrieval method works as follows :

STEP 1 - Visual Matching : Each query image feature $\mathbf{F}_j^q$ is matched to the full features dataset thanks to an efficient similarity search technique (see Section 3). It typically returns the approximate K-nearest neighbors according to the used metric d (i.e the K most similar pictures). When multiple matches occur for a given query image feature and a given retrieved record, we only keep the best match according to the feature distance. The visual matching step finally returns a set of candidate event records E_i, each being associated with M_i^q picture matches of the form (I_m^q, I_m^i).

STEP 2 - Stop List : Only the retrieved records with at least two image matches are kept for the next step, i.e

$$\{E_i \mid M_i^q \geq 2\}_{1 \leq i \leq N}$$

STEP 3 - Geo-temporal consistency : For each remaining record, we compute a geo-temporal consistency score by estimating a translation model between the query record and the retrieved ones. The resulting scores $S_q(E_i)$ are used to produce the final records ranking returned for query E_q. The translation model estimation is based on a robust regression and can be expressed as :

$$\widehat{\Delta}(E_q, E_i) = \arg\min_{\Delta} \sum_{m=1}^{M_i^q} \rho_\theta \left(\mathbf{P}_m^q - (\mathbf{P}_m^i + \Delta) \right) \tag{3.1}$$

where $\mathbf{P}_m^q$ and $\mathbf{P}_m^i$ are the geo-temporal coordinates of the m-th match (I_m^q, I_m^i). The cost function ρ_θ is typically a robust M-estimator allowing outliers to be rejected with a tolerance θ (in our experiments we used Tukey's robust estimator). The estimated translation parameter $\hat{\Delta}$ should be understood as the spatial and temporal offset required to register the query event record E_q with the retrieved

event record E_i. Once this parameter has been estimated, the final score of an event E_i is finally computed by counting the number of inliers, i.e the number of visual matches that respect the estimated translation model :

$$S_q(E_i) = \sum_{m=1}^{M_i^q} \left(\left\| \mathbf{P}_m^q - (\mathbf{P}_m^i + \widehat{\Delta}) \right\| \leq \theta \right) \tag{3.2}$$

where θ is a tolerance error parameter, typically the same as the one used during the estimation phase. In practice, we use a smooth counting operator to get a better dynamic on resulting scores. When we restrict ourselves to temporal metadata (as was done in the experiments), Equation 3.1 can be simplified to :

$$\hat{\delta}(E_q, E_i) = \arg\min_{\delta} \sum_{m=1}^{M_i^q} \rho_\theta \left(t_m^q - (t_m^i + \delta) \right) \tag{3.3}$$

where $\hat{\delta}$ represents the estimated temporal offset between E_q and E_i and θ is now a temporal tolerance error whose value is discussed in the experiments. Since δ is a single mono-dimensional parameter to be estimated, Equation 3.3 can be resolved efficiently by a brute force approach testing all possible solutions δ.
Final scores then become :

$$S_q(E_i) = \sum_{m=1}^{M_i^q} \left(\left| t_m^q - (t_m^i + \hat{\delta}) \right| \leq \theta \right) \tag{3.4}$$

STEP 4 - Prior constraints : Depending on the application context, major improvements in effectiveness might be obtained by adding prior constraints on the tolerated values for $\widehat{\Delta}$. Rejecting events with too large spatial and/or temporal offset from the query record is indeed a good way to reduce the probability of false alarms. In our experiments we study the impact of such a constraint on the estimated temporal offsets. Concretely, we reject from the result list all retrieved event records which have an estimated offset above a given threshold δ_{max} (regardless of the matching score $S_q(E_i)$).

3 Enabling scalability

To allow fast visual matching in large picture datasets, we implemented a distributed similarity search framework based on Multi-Probe Locality Sensitive Hashing [69, 53] and the MapReduce [28] programming model.

3.1 Multi-Probe LSH

To process the Nearest Neighbors search efficiently, we use an approximate similarity search structure, namely Multi-Probe Locality Sensitive Hashing (MP-LSH) [69, 53]. MP-LSH methods are built on the well-known LSH technique [24], but they intelligently probe multiple buckets that are likely to contain results. Such techniques have been proved to overcome the over-linear space cost drawback of common LSH while preserving a similar sub-linear time cost (with complexity $O(N^{\lambda})$).

Now, let $\mathcal{F}$ be the dataset of all visual features $\mathbf{F} \in \mathbb{R}^d$ (i.e. the one extracted from the pictures of the N event records E_i). Each feature $\mathbf{F}$ is hashed with a hash function $g : \mathbb{R}^d \rightarrow \mathbb{Z}^k$ such that :

$$g(\mathbf{F}) = (h_1(\mathbf{F}), ..., h_k(\mathbf{F})) \tag{3.5}$$

where individual hash functions h_j are drawn from a given locality sensitive hashing function family. In this work we used the following binary hash function family which is known to be sensitive to the inner product :

$$h(\mathbf{F}) = sgn\,(\mathbf{W}.\mathbf{F}) \tag{3.6}$$

where $\mathbf{W}$ is a random variable distributed according to $\mathcal{N}(0, \mathbf{I})$. The hash codes produced $\mathbf{g}_i = g(\mathbf{F}_i)$ are thus binary hash codes of size k.

At indexing time, each feature $\mathbf{F}_i$ is mapped into a single hash table $\mathbf{T}$ accor-

ding to its hash code value $\mathbf{g}_i$. As a result, we obtain a hash table of $\mathbf{N}_b$ buckets where $\mathbf{N}_b \leq 2^k$.

At query time, the query vector $\mathbf{F}_q$ is also mapped onto the hash table $\mathbf{T}$ according to its hash code value $\mathbf{g}_q$. The multi-probe algorithm then selects a set of $\mathbf{N}_p$ buckets $\{(\mathbf{b}_j)\}_{j=1..N_p}$ as candidates that may contain objects similar to the query according to :

$$d_h(\mathbf{g}_q, \mathbf{b}_j) < \delta_{MP} \tag{3.7}$$

where $\mathbf{d}_h$ is the hamming distance between two binary hash codes and δ_{MP} is the multi-probe parameter (i.e. a radius of hamming space).

A final step is then performed to filter the features contained in the selected buckets by computing their distance to the query and keeping the K Nearest Neighbors.

3.2 The MapReduce framework

MapReduce is a programming model introduced by Google to support distributed batch processing on large data sets. A MapReduce job splits the input dataset into independent chunks which are processed by the *map* tasks in a parallel manner. The framework sorts the outputs of the maps, which are then input to the *reduce* tasks. Chunks are processed based on key/value pairs. The *map* function computes a set of intermediate key/value pairs and, for each intermediate key, the *reduce* function iterates through the values that are associated with that key and outputs 0 or more values. The *map* and *Reduce* tasks scheduling is performed by the framework. In a distributed configuration, the framework assigns jobs to the nodes as slots become available. The number of *map* and *reduce* slots as well as chunk size can be specified for each job, depending on the cluster size. With such a granularity, large data sets processing can be distributed efficiently on commodity clusters.

3.3 Multi-Probe LSH in the MapReduce framework

The hash table T in the MapReduce framework is stored in a text file where each line corresponds to a single bucket. Each bucket is represented by a $< key, value >$ pair :

$$< \mathbf{b}, ((id(\mathbf{F}_1), \mathbf{F}_1), (id(\mathbf{F}_2), \mathbf{F}_2), \dots) > \tag{3.8}$$

where b is the hash code of the bucket and $id(\mathbf{F})$ the picture identifier associated to feature F.

In order to be processed by the MapReduce framework, the table T has to be divided into a set of splits. The number of splits is deduced by the MapReduce framework according to a set of input parameters as the number of available slots and the minimal input split size which is related to the file system block size. However, in order to be entirely processed by a mapper, a bucket cannot spill over different splits.

Since MapReduce is mainly dedicated to batch processing, setting up tasks could be expensive due to process creation and data transfer. Therefore, our implementation processes multiple queries at a time, typically sets of pictures belonging to the same records.

The hash codes of all query features are computed and passed to the *map* instances to be executed on the different slots. The number of *map* instances is computed by the MapReduce framework according to the number of input splits.

Each *map* process iterates over its assigned input split and for each query selects the candidate buckets that are likely to contain similar features according to Equ.3.7. It then computes the distance to each feature within the selected buckets. For each visited feature $\mathbf{F}_i$, the *map*function outputs a $< key, value >$

pair of the form :

$$< id(\mathbf{F}_q), (dist(\mathbf{F}_q, \mathbf{F}_i), id(\mathbf{F}_i)) > \tag{3.9}$$

where $dist(\mathbf{F}_q, \mathbf{F}_i)$ denotes the distance between $\mathbf{F}_q$ and $\mathbf{F}_i$.

For each query identifier $id(\mathbf{F}_q)$ the *reduce* instance sorts the set of emitted values for all *map* instances and filters the K-nearest neighbors.

Figure 3.3 gives the average response time per K-NN search according to the total number of queries batched within the same MapReduce job. It shows that the MapReduce framework becomes profitable from about 50 grouped queries. The average response time becomes almost constant for more than 400 grouped queries. In our experiments, the number of images per event record ranges from about 5 to 200. That means that using the MapReduce framework is still reasonable for the online processing of a single event record.
Finally, many MapReduce implementations materialize the entire output of each *map* before it can be consumed by the *reducer* in order to ensure that all *maps* successfully completed their tasks. In [22], Condel et al. propose a modified MapReduce architecture that allows data to be pipelined between operators. This extends the MapReduce programming model beyond batch processing, and can reduce completion times while improving system utilization for batch jobs as well.

4 Experiments

We evaluated our method on a *Flickr* image dataset using *last.fm* tags as real-world events ground truth. It was constructed from the corpus introduced by Troncy et al. [95] for the general evaluation of event-centric indexing approaches. This corpus mainly contains events and media descriptions and was originally created from three large public event directories (*last.fm*, *eventful* and *upcoming*). In our case, we only used it to define a set of Flickr images labeled with *last.fm*

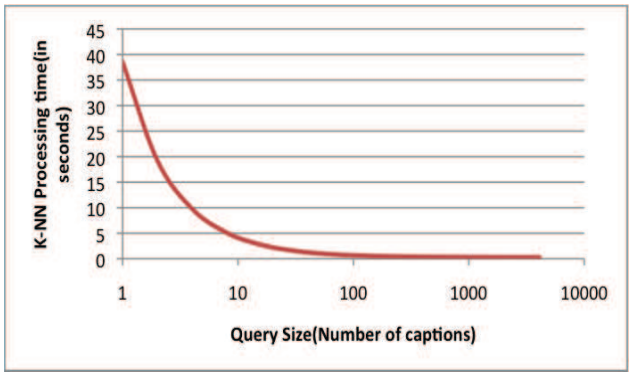


FIGURE 3.3 – Processing time per image according to query size

tags, i.e. unique identifiers of music events such as concerts, festivals, etc. The images themselves were not provided in the data and had to be crawled resulting in some missing images. Unfortunately, in this corpus, only a small fraction had geo-tags so that we evaluated our method using only temporal metadata. We used the EXIF *creation date* field of the pictures to generate the time metadata used in our method. Only about 50% of the crawled images had such a valid EXIF (others had empty or null date fields). In Table 3.1, we report the statistics on the original, crawled and filtered dataset. To gather the pictures in relevant event records, we used both the *last.fm* identifier and the *Flickr* author field provided with each picture. An event record is then defined as the set of pictures by a given author having the same LastFM label. Our final dataset contains $41,294$ event records related to $34,034$ distinct LastFM events.

TABLE 3.1 – Test dataset Vs Original dataset

	Total	Crawled	Filtered
photos	1 667 317	1637585	828902
users	23 060	22676	10257

4.1 Experimental settings

We used 6 global visual features to describe a picture's visual content (including HSV Histogram[34], Hough histogram[34], Fourier histogram[34], edge orientation histogram[34]). Each feature was $L2$-normalized and hashed into a 1024 bits hash code using the same hash function as the one used to construct the hash table (see Equ.3.6). The 6 hash codes were then concatenated into a single hash code of 6144 bits. We used the Hamming distance on these hash code as visual similarity. From the full set of $41,294$ event records in the dataset, the only queries we kept were the records being tagged with *last.fm* events and having at least 7 records in the dataset. We finally got 172 query records E_q. This procedure was motivated by the fact that a very large fraction of events were represented by only one record and therefore not suitable for experiments.

In all the experiments, we used a leave-one-out evaluation procedure and measured performances with 2 evaluation metrics : Mean Average Precision (MAP) and Classification Rate (CR). MAP is used in most information retrieval evaluations and measures the ability of our method to retrieve all the records related to the same event as the query event. Classification rate is obtained by using our method as a nearest neighbors classifier. The number of occurrences of retrieved events is computed from the top 10 returned records and we keep the event with the maximum score as the best prediction. It measures the ability of our method to automatically label some unknown query event record. We extend this measure to the case of multiple labels suggestion. In addition to the best retrieved event we also return the following events by decreasing scores (*i.e.* decreasing number of occurrences found within the top-10 returned records). In this case, the success rate is measured by the percentage of query records where the correct event was retrieved among all suggested event tags. It measures the performance of our method in the context of tags suggestion rather than automatic annotation.

Finally, we used the Hadoop [1] MapReduce implementation on a 5-node cluster. Nodes are equipped with Intel Xeon X5560 CPUs as well as 48Gb of RAM.

4.2 Results

Parameters discussion

In Figure 3.4, we report the mean average precision for varying values of the θ parameter (Eq. 3.3) and different numbers of K-nearest neighbors used during the visual matching step. The results show that MAP values are at their optimal for $\theta \in [300, 1800]$ seconds. This optimal error tolerance value is coherent with the nature of the events in the *last.fm* corpus. Picture records of concerts indeed usually range from one to several hours. On the other hand, below 5 minutes, real-world concert scenes are too ambiguous to be discriminated by their visual content (or at least with the global visual features used in this study). In what follows, we fix θ to 1800 as an optimal value.

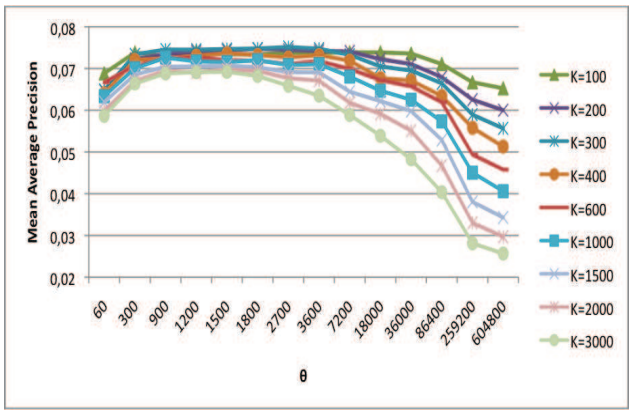


FIGURE 3.4 – Influence of temporal error to tolerance θ

We now study the impact of adding a prior constraint δ_{max} on the estimated temporal offsets $\hat{\delta}$. Most events in *last.fm* dataset being music concerts, it is un-

1. http ://hadoop.apache.org/mapreduce/

likely that the temporal offset between two records would reach high values. We therefore study the impact of rejecting all retrieved records having a temporal offset higher than δ_{max}. Figure 3.5 displays the new MAP curves for varying values of δ_{max}. It shows that the mean average precision can be consistently improved from about 0.08 without any constraint to 0.18. The optimal value for δ_{max} is about 86,400 seconds which is exactly 1 day. That means that the records of a single real-word event might have a temporal offset of up to 1 day. The interpretation is that the EXIF *creation date* field is noisy due to the different reference times of the devices used (users from different countries, default device settings, etc.). It is worth noting that our method is by its very nature robust to such temporal offsets since we mainly consider temporal coherence rather than absolute time matching. On the other hand, rejecting records with temporal offsets higher than 1 day allows many visual false positives to be rejected.

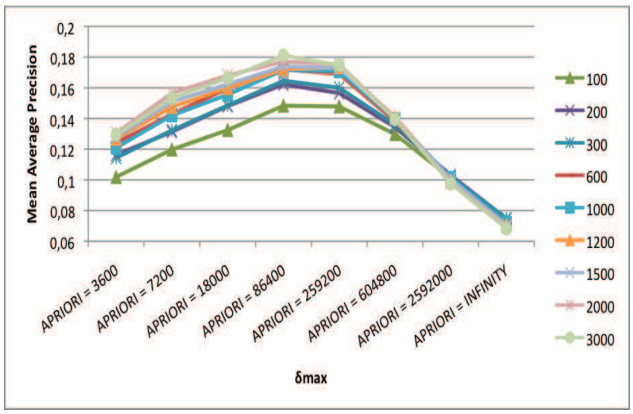


FIGURE 3.5 – Influence of temporal offset thresholding (δ_{max}) on MAP

Figure 3.6 displays the results of the same experiment but for the classification rate (using a 10-NN classifier on retrieved records) rather than the mean average precision. This evaluates the ability of our method to automatically annotate a query event record rather than its ability to retrieve all records in the dataset. Here again the optimal classification rates are obtained when δ_{max}=1 day. Fur-

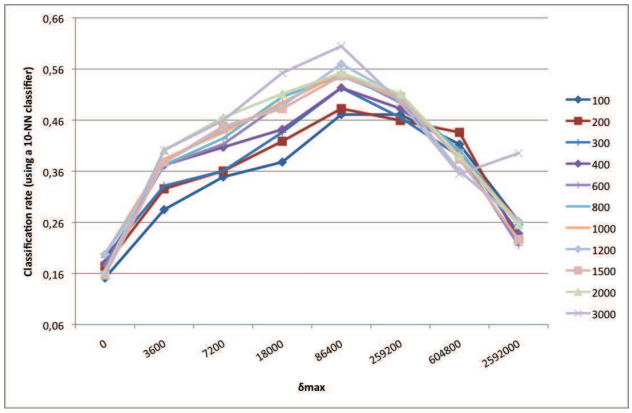
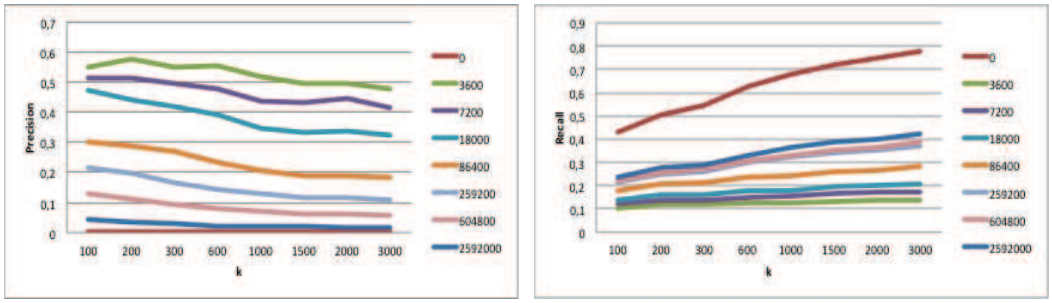


FIGURE 3.6 – Influence of temporal offset thresholding (δ_{max}) classification rates

thermore, we see that the classification rate always increases with the number K of closest visual matches (returned for each query image). The interpretation is that increasing K improves recall without degrading precision too much thanks to the selectivity of our temporal consistency re-ranking step. We verified this by studying the recall and the precision independently.

Figure 3.7 displays both precision and recall for increasing values of K. The results confirm the above conclusion that increasing values of K improves recall without compromising much of the precision. This shows the ability of our temporal consistency re-ranking step to efficiently surface relevant records.



(a) Precision (b) Recall

FIGURE 3.7 – Precision and recall for increasing values of k

Event suggestion in the MapReduce framework

All the previous experiments were made using an exhaustive search for the k-NN search. In this section, we evaluate the performance of our full framework using MapReduce and the Multi-Probe LSH. As parameters, we used the optimal values discussed in the previous section (i.e. δ_{max}=86.400, K=3000 and θ=1800).

Table 3.2 displays the class rates using an exhaustive search as seen in the previous section as well as class rates using a Multi-Probe LSH-based similarity search for different values of δ_{MP}.

TABLE 3.2 – Suggestion rates

# of suggested events tag	1	2	3	4	5	10
Exhaustive	0.60	0.66	0.69	0.71	0.72	0.73
MP-Delta 0	0.39	0.48	0.50	0.51	0.52	0.54
MP-Delta 1	0.45	0.55	0.57	0.58	0.59	0.63
MP-Delta 2	0.48	0.59	0.61	0.65	0.66	0.69
MP-Delta 4	0.51	0.61	0.63	0.67	0.67	0.70
MP-Delta 8	**0.61**	**0.67**	**0.70**	**0.72**	**0.72**	**0.74**
MP-Delta 16	0.59	0.66	0.69	0.71	0.72	0.73

As one might expect, all class rates values increase accordingly with the number of probes (i.e increasing δ_{MP} values) to surprisingly perform better than the exhaustive search for δ_{MP}=8. Overall, in the best case, our method is able to suggest the correct event tag over 5 suggestions with a 72% success rate. Such performances are clearly acceptable from an application point of view.

Figure 3.8 displays the average search time per query for both distributed and centralized search. We compare the K-NN processing time per image for a centralized setting (number of *map* slots = 1) to the processing time in a distributed scheme (20 *map* slots available on the network) for both exact and approximate similarity search. Although the multi-probe might reduce the effectiveness of our

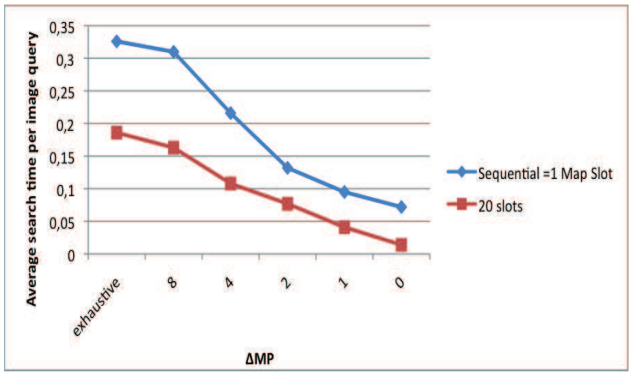


FIGURE 3.8 – K-NN search time per image ($k = 4000$)

method, it might significantly reduce the search time by an order of magnitude.

4.3 Dicussion

On event matching

From an application point of view, the identification rates achieved are clearly acceptable for event search and identification. Yet, we believe that the identification rate can be further improved with more information such as geotags and user generated annotations (people tags, contextual annotations, etc).

Obviously, linking different records of the same event would also enable the discovery of event related documents in social media. Furthermore, this would enable the discovery of more complex relationships and patterns between the event. In order to handle such date sets effectively, our method should hence, scale accordingly.

On scalability

The performance gain achieved by Multi-Probe LSH over exhaustive search is nonetheless still less than the one obtained in usual centralized settings. First, in MapReduce approaches, probing multiple buckets generates more network ove-

rhead in addition to data transfers across the network. The second reason is due to bucket occupation. In fact, imbalanced buckets generate imbalanced map chunks leading to disproportionate map execution times.

While recent efforts have overcome some of the Hadoop MapReduce Framework limitations and more recently the Apache Hadoop NextGen MapReduce (YARN) [2], computing the K-Nearest Neighbors for each image is time and space consuming even if they are processed in a parallel and/or distributed manner since it often implies distance computation or approximation.

In Chapter 4, we address the problem of computing the images K-Nearest Neighbors in a fully decentralized, scalable and space-adaptive manner.

5 Conclusion

In this chapter we presented a new visual-based method for retrieving events in photo collections, that might also be used for event tag suggestion or annotation. Our method proved to be robust to temporal offsets since we mainly rely on temporal coherence rather than absolute time matching. As one result, we are able to suggest the correct event tag with a success rate of at least 60% and even 72% if we allow multiple suggestions.

The proposed method is scalable, since it relies on efficient approximate similarity search techniques based on the MapReduce framework. We also investigated multi-probe techniques trading accuracy for efficiency, which might lead to a loss of 8.3% class rate compared to a gain of 58.6% in processing time.

2. http ://hadoop.apache.org/docs/r0.23.0/hadoop-yarn/hadoop-yarn-site/YARN.html

Chapitre 4

A distributed Framework for k-NN Graphs construction

Efficiently constructing the K-Nearest Neighbor Graph (K-NNG) of large and high dimensional datasets is crucial for many applications with feature-rich objects, such as images or other multimedia content. In this chapter we investigate the use of high dimensional hashing methods for efficiently approximating the K-NNG in distributed environments. We first discuss the importance of balancing issues on the performance of such approaches and show why the baseline approach using Locality Sensitive Hashing does not perform well. Our new KNN-join method is based on RMMH, a recently introduced hash function family based on randomly trained classifiers. We show that the resulting hash tables are much more balanced and that the number of resulting collisions can be greatly reduced without degrading quality. We further improve the load balancing of our distributed approach by designing a parallelized local join algorithm. We show that our method outperforms state-of-the-art methods in centralized settings and that it is efficiently scalable given its inherently distributed design. Finally, we present a distributed implementation of our method using a MapReduce framework and evaluate its performance on a large dataset.

1 Problem Statement

Given a set $\mathcal{X}$ of N objects, the K-Nearest Neighbor Graph consists of the vertex set $\mathcal{X}$ and the set of edges connecting each object from $\mathcal{X}$ to its K most similar objects in $\mathcal{X}$ under a given metric or similarity measure. Efficiently constructing the K-NNG of large datasets is crucial for many applications involving feature-rich objects, such as images, text documents or other multimedia content. Examples include query suggestion in web search engines [85], collaborative filtering [1], visual objects discovery [80] and event detection in multimedia User Generated Contents. The K-NNG is also a key data structure for many established methods in data mining [12], machine learning [10] and manifold learning [103]. Overall, efficient K-NNG construction methods would extend a large pool of existing graph and network analysis methods to large datasets without an explicit graph structure.

In this chapter we investigate the use of high dimensional hashing methods for efficiently approximating the K-NNG, notably in distributed environments. A decade after the first LSH [37], hashing methods have indeed attracted increasing interest for efficiently solving Nearest Neighbors problems in high-dimensional feature spaces. Embedding high-dimensional feature spaces in very compact hash codes makes it possible to scale up many similarity search applications (from 10 to 1000 times larger datasets) [38, 55, 101]. One advantage of hashing methods over trees or other structures is that they simultaneously allow efficient indexing and data compression. Hash codes can indeed be used to gather features into buckets but also to approximate exact similarity measures by efficient hash code comparisons (typically a hamming distance on binary codes). Memory usage and processing costs can therefore be drastically reduced.

Unfortunately, recent studies [107, 31] have shown that the performances of usual hashing-based methods are not as good as expected when constructing the full K-NNG (rather than only considering individual top-K queries). Recently, Dong et al.[31] even show that LSH and other hashing scheme can be outperformed by a radically different strategy purely based on query expansion operations

[31], without relying on any indexing structure or partitioning method. Our work provides evidence to support hashing based methods by showing that such observations might be mitigated when moving to more recent hash function families.

Our new KNN-join method is notably based on RMMH [55], a recent hash function family based on randomly trained classifiers. In this chapter, we discuss the importance of balancing issues on the performance of hashing-based similarity joins and show why the baseline approach using Locality Sensitive Hashing (LSH) and collisions frequencies does not perform well (Section 2). We then introduce our new K-NNG method based on RMMH (Section 3). To further improve load balancing in distributed environments, we finally propose a distributed local join algorithm and describe its implementation within the MapReduce framework.

2 Hashing-based K-NNG construction

2.1 Notations

Let us first introduce some notations. We consider a dataset $\mathcal{X}$ of N feature vectors $\mathbf{x}$ lying in a Hilbert space $\mathbb{X}$. For any two points $\mathbf{x}, \mathbf{y} \in \mathcal{X}$, we denote as $\kappa : \mathbb{X}^2 \to \mathbb{R}$ a symmetric kernel function satisfying Mercer's theorem, so that κ can be expressed as an inner product in some unknown Hilbert space through a mapping function Φ such that $\kappa(\mathbf{x}, \mathbf{y}) = \Phi(\mathbf{x}).\Phi(\mathbf{y})$.

If we denote by $\mathcal{N}_K(\mathbf{x})$ the set of the K nearest neighbors of $\mathbf{x}$ in $\mathcal{X}$ according to κ, then the K-Nearest Neighbor Graph on $\mathcal{X}$ is a directed graph $\mathbf{G}_K(\mathcal{X}, E)$ connecting each element to its K Nearest Neighbors, thus :

$$E = \{(\mathbf{u}, \mathbf{v}), \mathbf{u} \in \mathcal{X}, \mathbf{v} \in \mathcal{N}_K(\mathbf{u})\}$$

We generally denote by $\mathcal{H}$, a family of binary hash functions $h : \mathbb{X} \to \{-1, 1\}$. If we consider hash function families based on random hyperplanes we have :

$$h(\mathbf{x}) = sgn\,(\mathbf{w}.\mathbf{x} + b)$$

where $\mathbf{w} \in \mathbb{X}$ is a random variable distributed according to p_w and b is a scalar random variable distributed according to p_b. When working in the Euclidean space $\mathbb{X} = \mathbb{R}^d$ and choosing $p_w = \mathcal{N}(0, \mathbf{I})$ and $b = 0$, we get the popular LSH function family sensitive to the inner product [15, 60]. In this case, for any two points $\mathbf{q}, \mathbf{v} \in \mathbb{R}^d$ we have :

$$Pr[h(\mathbf{q}) = h(\mathbf{v})] = 1 - \frac{1}{\pi} cos^{-1}\left(\frac{\mathbf{q}.\mathbf{v}}{\|\mathbf{q}\| \, \|\mathbf{v}\|}\right)$$

Thus, the collision probability of any two items increases with their inner product ($\kappa(\mathbf{q}, \mathbf{v}) = \mathbf{q}.\mathbf{v}$). More generally, any LSH function family has the property :

$$Pr[h(\mathbf{q}) = h(\mathbf{v})] = f\left(\kappa(\mathbf{q}, \mathbf{v})\right) \tag{4.1}$$

where $f(\kappa)$ is the *sensitivity function*, increasing with κ.

2.2 LSH based K-NNG approximation

Let us consider L hash tables, each constructed from the concatenation of p hash functions built from an LSH family $\mathcal{H}$. The collision probability of any two items $\mathbf{q}, \mathbf{v}$ in *one* table is :

$$Pr[\mathbf{h}(\mathbf{q}) = \mathbf{h}(\mathbf{v})] = [f\left(\kappa(\mathbf{q}, \mathbf{v})\right)]^p \tag{4.2}$$

where $\mathbf{h}(\mathbf{q})$ and $\mathbf{h}(\mathbf{v})$ denote p-length binary hash codes.
The total number of collisions in *any set* of L hash tables, denoted as $n_{\mathbf{q},\mathbf{v}}$, is a random variable distributed according to a binomial distribution with parameters L (the number of experiments) and f^p (the probability of success of each Bernouilli experiment). The *expected* number of collisions in L hash tables is therefore :

$$\bar{n}_{\mathbf{q},\mathbf{v}} = E[n_{\mathbf{q},\mathbf{v}}] = L.\left[f\left(\kappa(\mathbf{q}, \mathbf{v})\right)\right]^p$$

The *empirical* number of collisions in L tables, denoted as $\hat{n}_{\mathbf{q},\mathbf{v}}$, can be seen as an estimator of this value. And since the sensitivity function f is supposed to be an increasing function with κ, it is easy to show that :

$$\kappa(\mathbf{q},\mathbf{v}_1) < \kappa(\mathbf{q},\mathbf{v}_2) \Leftrightarrow E[\hat{n}_{\mathbf{q},\mathbf{v}_1}] < E[\hat{n}_{\mathbf{q},\mathbf{v}_2}] \tag{4.3}$$

The top-K neighbors of any item $\mathbf{x} \in \mathcal{X}$ according to κ can therefore be approximated by the top-K items ranked according to their collision frequency with $\mathbf{x}$ (as suggested in [64]). Consequently the whole K-NNG on $\mathcal{X}$ can be approximated by simply counting the number of collisions of item pairs, without any distance computation.

More formally, we define the *hashing based approximation* of a K-NNG $\mathbf{G}_K(\mathcal{X}, E)$, as a new directed graph $\hat{\mathbf{G}}_K(\mathcal{X}, \hat{E})$ where $\hat{E}$ is a set of edges connecting any item $\mathbf{x}$ to its K most frequently colliding items in the L hash tables. In practice, since the number of collisions is a discrete variable, more than K items might have the same number of collisions and have to be kept in the graph produced. The hash-based approximation of a K-NNG should therefore rather be seen as a filtering step of the all-pairs graph. A brute-force refinement step can be applied on $\hat{\mathbf{G}}_K(\mathcal{X}, \hat{E})$ to get a more accurate approximation during a second stage.

2.3 Balancing issues of LSH-based K-NNG

The LSH-based K-NNG approximation is very attractive in the sense that it does not require any kernel (or metric) computation. It simply requires building L hash tables and post-processing all collisions occurring in these tables. Unfortunately, balancing issues strongly affect the performance of this scheme in practice. The cost of the method is, in fact, mainly determined by the total number of collisions in all hash tables, i.e.

$$T_{\mathcal{H}}(\mathcal{X}, L, p) = \sum_{l=1}^{L} \sum_{b=1}^{2^p} \frac{n_{l,b}.(n_{l,b}-1)}{2} \tag{4.4}$$

where $n_{l,b}$ is the number of items in the b-th bucket of the l-th table. For an ideally balanced hash function and $p \sim log_2(N)$, the cost complexity would be $O(L.N)$. But for highly unbalanced hash functions, the cost complexity tends rather to be $O(L.N^2)$ because the most filled buckets concentrate a large fraction of the whole dataset (i.e. $n_{l,b} = \alpha N$). To illustrate the potential impact of LSH balancing issues, Table 4.1 provide some real statistics computed on one of the datasets used in our experiments (see section 5), compared to a perfectly balanced hash function (L=128, p=16). It shows that the number of collisions to be processed is 3 orders of magnitude greater than the perfectly balanced hash function, resulting in an intensive computing cost. The poor balancing performance is confirmed by a very bad Gini coefficient and a low entropy. Overall, several authors have confirmed that LSH-based methods for approximating K-NN Graphs are not very efficient [31, 107]. Balancing however was not identified as being critical to improve the efficiency of hash-based K-NNG approximations.

Hash function	Perfect	LSH
Nb of collisions $T_{\mathcal{H}}(\mathcal{X}, L, p)$	$4.82 * 10^6$	$7.57 * 10^9$
Entropy	16	7.58
Gini coeff.	0	0.94
Max bucket size	12	100751
Nb of non empty buckets	65 536	11 070

TABLE 4.1 – Balancing statistics of LSH vs. perfectly balanced hash function

3 Proposed method

We now describe our K-NNG approximation method. It can be used either as a filtering step (combined with a brute-force refinement step applied afterwards), or as a direct approximation of the graph, depending on the quality of the application required. The method holds for centralized settings as well as for distributed or parallelized settings, as discussed below.

3.1 Random Maximum Margin Hashing

Rather than using classical LSH functions, our method is based on Random Maximum Margin Hashing (RMMH, [55]), an original hash function family introduced recently and one that is suitable for any kernelized space (including the classical inner product). In addition to its nice embedding properties, the main strength of RMMH for our problem is its load balancing capabilities. The claim of this method is actually that the lack of independence between hash functions is the main issue affecting the performance of data dependent hashing methods compared to data independent ones. Indeed, the basic requirement of any hashing method is that the hash function provide a **uniform** distribution of hash values, or at least one that is as uniform as possible. Non-uniform distributions increase the overall expected number of collisions and therefore the cost of resolving them. The uniformity constraint should therefore not be relaxed too much, even if we aim to maximize the collision probability of close points.

The main originality of RMMH is to **train** purely **random splits** of the data, regardless of the closeness of the training samples (i.e. without any supervision). The authors showed that such a *data scattering* approach makes it possible to generate consistently more independent hash functions than other data-dependent hashing functions. Moreover, the use of large margin classifiers allows good generalization performances to be maintained.

Concretely, the method works by learning a set of *randomly trained classifiers* from a small fraction of the dataset. For each hash function, M training points are selected at random from $\mathbf{X}$ and are then **randomly labeled** (half of the points with -1 and the other half with 1). If we denote as x_j^+ the resulting $\frac{M}{2}$ positive training samples and as $\mathbf{x}_j^-$ the $\frac{M}{2}$ negative training samples, each hash function is then computed by training a binary classifier $h_\theta(\mathbf{x})$ such that :

$$h(\mathbf{x}) = \arg\max_{h_\theta} \sum_{j=1}^{\frac{M}{2}} h_\theta(\mathbf{x}_j^+) - h_\theta(\mathbf{x}_j^-) \tag{4.5}$$

Using a Support Vector Machine (SVM) as a binary classifier, we get :

$$h(\mathbf{x}) = sgn\left(\sum_{i=1}^{m} \alpha_i^* \kappa(\mathbf{x}_i^*, \mathbf{x}) + b_m\right) \tag{4.6}$$

where $\mathbf{x}_i^*$ are the m support vectors selected by the SVM ($\mathbf{x}_i^* \in \left\{\mathbf{x}_j^+, \mathbf{x}_j^-\right\}$).

In the linear case ($\kappa = inner\ product$), this simplifies to :

$$h(\mathbf{x}) = sgn\left(\mathbf{w}.\mathbf{x} + b\right) \tag{4.7}$$

with $\mathbf{w} = \sum_{i=1}^{m} \alpha_i^* \mathbf{x}_i^*$ and the hash function is much faster to compute.

3.2 RMMH-based K-NNG approximation

Our K-NNG approximation algorithm now works in the following way :

— **STEP1 - Hash tables construction** : For each item $\mathbf{x} \in \mathcal{X}$, we compute L p-length binary hash codes $\mathbf{h}(\mathbf{x})$ (using $L.p$ distinct RMMH functions) and insert them in L distinct hash tables.

— **STEP2 - Local joins** : Non-empty buckets of each hash table are processed independently by a *local join* algorithm. For each non-empty bucket b, the local join algorithm generates the $n_c = \frac{n_b.(n_b-1)}{2}$ possible pairs of items computed from the n_b items contained in the bucket. Notice that this algorithm ensures that buckets can be processed separately and therefore facilitate the distribution of our method. Each emitted pair is simply coded by a pair of integer identifiers (id_i, id_j) such that $id_i < id_j$ (as a coding convention) with $0 < i < n_b - 1$ and $0 < i.n_b + j < n_c - 1$.

— **STEP3 - Reduction** : All pairs (id_i, id_j) are mapped onto an *accumulator* in order to compute the occurrence of each pair (within the $T_{\mathcal{H}}$ emitted pairs). Notice that the occurrence matrix produced does not depend on the mapping sequence so that each pair can be inserted independently from the other ones, at any time during the algorithm. This ensures that this reduction step can be easily distributed.

— **STEP4 - Filtering** : Once the full occurrence matrix has been computed, it is filtered in order to keep only the most similar items to each candidate item and compute our approximate graph $\hat{\mathbf{G}}_K(\mathcal{X}, \hat{E})$. This is done by scanning each line of the occurrence matrix and maintaining a priority queue according to the number of occurrences of each pair. Since the number of occurrences is a *discrete* priority value, all items having the same frequency are pulled together from the queue, so we finally get more than K most similar items for each candidate item. Notice that each line of the matrix can be processed independently from the other ones. This ensures that this filtering step can be easily distributed.

We recall here that the hashing-based K-NNG produced by the above algorithms could still be refined by a brute-force algorithm applied on the remaining pairs. We use such a refinement step in some of our experiments so as to make comparisons possible with the state-of-the-art results of [31].

3.3 Split local joins

Although local joins can be easily distributed, large buckets still affect the overall performance. The local join algorithm has quadratic space complexity ($n_c = \frac{n_b.(n_b-1)}{2} = O(n_b^2)$) and is therefore, likely to raise memory exceptions as well as expensive swapping phases. Moreover, we want our distributed framework to support a wide variety of hash functions, even those with lower balancing capabilities such as LSH. In the following, we extend the local join algorithm to process large buckets in parallel and/or distributed architectures with guarantees on the runtime and memory occupation.

In practice, if the number of collisions generated by a given local join exceeds a fixed threshold (i.e. $n_c > c_{max}$), then the local join is split into $n_s = \lceil \frac{n_c}{c_{max}} \rceil$ *sub-joins*, each being in charge of at most c_{max} collisions.

Algorithm 1 shows the pseudo-code of the basic local join.

Since the number of generated pairs at iteration k is $n - k$, the number of

Algorithm 1 *Local Join*

Require: bucket $b = \{id_i\}_{1<i<n_b}, start, end.$
Ensure: local collisions set C.

1: $C \leftarrow \emptyset$
2: **for** $i \leftarrow start, \ldots, end$ **do**
3: **for** $j \leftarrow i+1 \ldots, n_b$ **do**
4: $C \leftarrow C \cup (b_i, b_j)$
5: **end for**
6: **end for**

generated pairs of an s-length iteration block starting at the i^{th} iteration of the external loop is :

$$\sum_{k=0}^{s-1}(n_b - i - k) = \frac{1}{2} * s^2 + (n_b - i - \frac{1}{2}) * s \tag{4.8}$$

which must be less than or equal to c_{max}. The thus defined inequality has two roots of opposite signs s_1 and s_2 $(s_1 > s_2)$; we require s to be equal to $\lfloor s_1 \rfloor$ as long as $s + i$ remains less than or equal to s, $n_b - i$ otherwise.

Algorithm 2 gives the pseudo-code for the enhanced local-join with the *split* strategy. It first computes the starting iteration of each iteration block (lines 3-10), local joins are then spawned concurrently across processing units (lines 11-13).

3.4 MapReduce Implementation

As explained in the previous section, all the steps of our hashing-based K-NNG approximation framework can be easily distributed. In this work, we implemented it under the Hadoop *MapReduce* framework [27]. This is probably not the most efficient implementation, but it is highly scalable and easily deployable into large computing clouds. A first *MapReduce* job performs the hash tables construction **STEP1** and then, a second MapReduce job computes **STEP2** and **STEP3** (using the split local join strategy). **STEP4** was not implemented under MapReduce within our experiment but this could be easily done by using the occurrence matrix line numbers as input keys to a third job.

Algorithm 2 *Distributed and/or parallel Local Join*

Require: bucket $b = \{id_i\}_{1<i<n_b}$, capacity c_{max}
Ensure: distributed collisions set

1: $l \leftarrow \emptyset$ *//starting iterations list*
2: $k \leftarrow 1$
3: **while** $k < n_b$ **do**
4: $\quad s_1 \leftarrow \lfloor \frac{1}{2} - n_b + k + \sqrt{(n_b - k - \frac{1}{2})^2 + 2.c_{max}} \rfloor$
5: $\quad$ **if** $s_1 > n_b - k$ **then**
6: $\quad\quad s_1 \leftarrow n_b - k$
7: $\quad$ **end if**
8: $\quad l \leftarrow l \cup s_1$
9: $\quad k \leftarrow k + s_1$
10: **end while**
11: **for** $i \leftarrow 1 \ldots, |l| - 1$ **do**
12: $\quad Local\ Join(b, l[i], l[i+1] - 1)$
13: **end for**

Hash table construction (STEP1)

The first MapReduce job splits the input dataset $\mathcal{X}$ into independent chunks of equal sizes to be processed in parallel. A mapper iterates over the set of its assigned object features and computes $L.p$ hash values for each feature according to Equation 4.6. Hash values are concatenated into L p-length hash codes corresponding to L bucket identifiers for the L hash tables). Each hash code is then emitted along with the table identifier (*intermediate key*) and the associated feature identifier (intermediate value).
The Reduce function merges all the emitted identifiers for a particular intermediate key (*i.e.* bucket identifier within a specific table). The resulting buckets are provided as input to the second MapReduce job.

Occurrence matrix computation (STEP2 & 3)

The second job processes buckets separately. The map function generates all possible pairs of identifiers of the processed bucket and issues each pair (*interme-*

diate key), possibly with a null intermediate value. The reduce function counts the number of intermediate values for each issued pair. For efficiency reasons, map outputs are combined locally before being sent to the reducer. This requires intermediate values to store the cumulated pair occurrences. With such an optimization, the mapper issues each pair along with its initial occurrence. Combine and reduce functions simply sum the intermediate values for each issued pair.

4 Experimental setup

This section provides details about the experimental setup, including datasets, performance measures, default parameters and system environment. Experimental results are reported in Section 5.

4.1 Datasets & Baselines

Our method was evaluated on 3 datasets of different dimensions and sizes :

Shape : a set of 544-dimensional feature vectors extracted from 28775 3D polygonal models from various sources.
Audio : a set of 54387 192-dimensional feature vectors extracted from the DARPA TIMIT collection.
Flickr : we use the same dataset as described in section 4.

All feature vectors were L_2 normalized and compressed into 3072-dimensional binary hash codes using RMMH and LSH. Table 4.2 summarizes the salient information of these datasets.

The shape and audio datasets were first used in [32] to evaluate the LSH method and more recently in [31] to evaluate the NN-*Descent* algorithm against the Recursive Lanczos Bisection[17] and LSH. We rely on these datasets to evaluate our method against the NN-*Descent* method [31] (which outperforms previous approximate KNNG methods).

Finally, we use the Flickr dataset to study in more detail the performances of

our method in the context of a larger dataset (in size and dimensionality), related to the context of this dissertation (*i.e.* event mining).

Datasets	# Objects	Dimension
Shape	28 775	544
Audio	54 387	192
Flcikr	828 902	793

TABLE 4.2 – Dataset summary

4.2 Performance measures

We use *recall* and *precision* to measure the **accuracy** of our approximate KNN Graphs against the exact KNN Graphs. The exact K-NN Graphs were computed on each dataset using a brute-force exhaustive search to find the K-NN of each node. The default K is fixed to 100. The default similarity measure between feature vectors is the inner product. Note that, since all features are L_2-normalized, the inner product K-NNG is equivalent to the Euclidean distance K-NNG. The *recall* of an approximate K-NNG is computed as the number of correct Nearest Neighbors retrieved, divided by the number of edges in the exact K-NNG. Similarly, we define the *precision* of an approximate K-NNG as the number of exact Nearest Neighbors retrieved, divided by the total number of edges in the approximate K-NNG.

The **efficiency** of our method is evaluated with the following metrics :

— Number of generated pairs : is used as an architecture-independent measure of the cost of our method to study the impact of the different parameters and hash functions used.

— Gini coefficient : is used to measure the load balancing of the hash tables used. Low *Gini* coefficients reflect good bucket balancing, while high values reveal large disparities in feature distribution. We should mention here that a null value reflects a uniform distribution in the hash space. In [81], the

authors show that the *Gini* coefficient is the most appropriate statistical metric for measuring load balancing fairness. For a better understanding of the impact of unfair feature balancing on the running time, we also report some statistics on the average maximum bucket size and the average number of non-empty buckets.

— Scan-Rate : is used as an architecture and method-independent measure of the filtering capabilities of approximate KNNG construction methods. It is defined in [31] as the ratio of the number of item pairs processed by the algorithm to the total number of possible pairs (i.e $\frac{N(N-1)}{2}$).

— CPU Time : is used to compare the overall efficiency of our method against the NN-*descent* method.

— Min, Max and Average Map running times are used to evaluate the performances that can be achieved on large clusters.

4.3 System environment

We implemented our approach on the Hadoop [1] MapReduce framework.

MapReduce-based experiments were conducted on a 6-node cluster, each node being equipped with four 2.8 *Ghz* Quad Core Intel Xeon CPU and 48 *Gbytes* of memory. The number of configured *map*, respectively *reduce* slots is hardware-dependent and is limited by the amount of available memory as well as the number of supported parallel threads per node. In order to avoid expensive context switches and memory swaps, we require each node to host at most 8 *map* slots and 3 *reduces* in parallel.

The NN-*Descent* code is the same as in [31] and was provided by the authors. It is an openMp parallelized implementation and runs only in centralized settings. To allow fair comparison, we used an openMp-based centralized version of our code rather than the MapReduce implementation. It iteratively performs steps 1 to 4 (Section 3.1) and finally applies a brute-force refinement step on the remaining pairs. Centralized experiments were conducted on an $X5675$ 3.06 Ghz processor

1. http ://hadoop.apache.org/mapreduce/

server with 96 *Gbytes* of memory.

5 Experimental results

We first evaluate the impact of the hash functions used on load distributions (Section 5.1). We then evaluate the overall performance of our method in centralized settings (Section 5.2) and compare it against the NN-*Descent* algorithm (Section 5.2).

The last part serves to validate our method in the MapReduce framework(Section 5.3).

5.1 Hash functions evaluation

We first evaluate the ability of RMMH to produce fair load distributions in the hash tables. This was not addressed in the original work of Joly et al. [55].

In Figure 4.1, we report the Gini coefficient for different values of M, i.e. the main parameter of RMMH (Section 3.1). The plots show that hash tables produced by RMMH quickly converge to fair load balancing when M increases. Gini coefficients below 0.6 are, for instance, known to be a strong indicator of a fair load distribution [44]. As a proof of concept, very high values of M even provide near-perfect balancing. As we will see later, such values are not usable in practice since too much scattering of the data also degrades the quality of the approximate graph generated. The parameter M is actually aimed at tuning the compromise between hash functions independence and the generalization capabilities of the method [55].

In Table 4.3, we report some statistics for LSH based hashing. Although LSH achieves correct balancing on the Shape and Audio datasets, it performs consistently worse on the Flickr dataset. For typical values of M greater than 15, RMMH

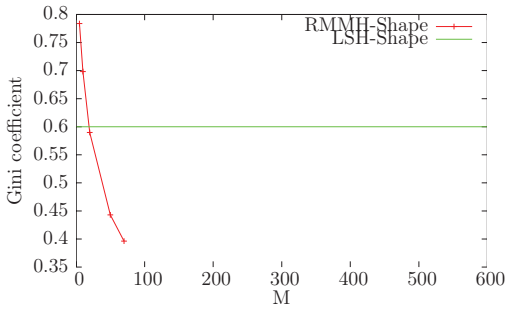


FIGURE 4.1 – Gini coefficient - RMMH-Based Hashing

outperforms LSH on the 3 datasets.

	Shape	Audio	Flickr
Gini	0.60	0.63	0.94
# non-empty buckets	6656	14917	11071
Avg. Max. bucket size	594	468.039	100751

TABLE 4.3 – Bucket Balancing Statistics - LSH-Based Hashing

This can be further verified in Figures 4.2 and 4.3 as the maximal bucket size per dataset decreases inversely to the number of non-empty buckets.

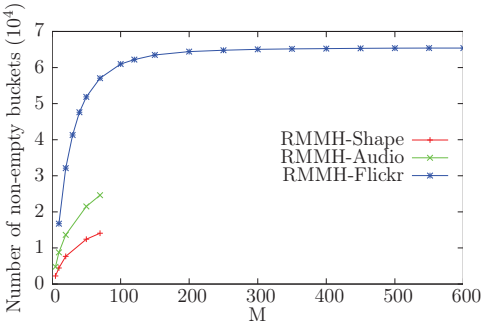


FIGURE 4.2 – # of non empty buckets - RMMH-Based Hashing

Figure 4.4 plots the number of collisions to be processed for increasing values of L and different hash functions. The results show that the RMMH based approach generates up to two orders of magnitude fewer collisions than the LSH-based ap-

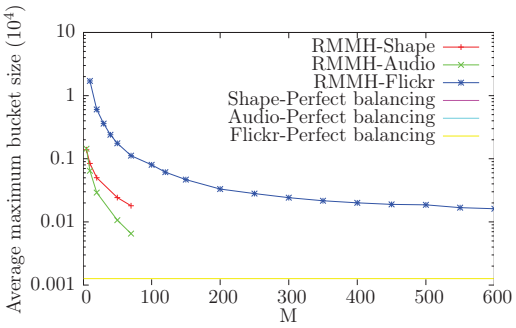


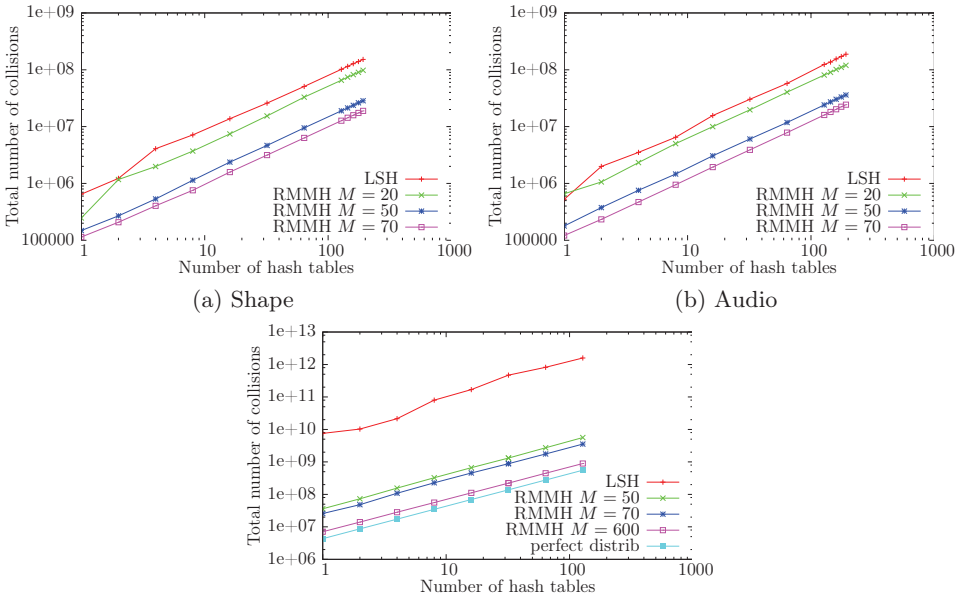
FIGURE 4.3 – Average maximum bucket size - RMMH-Based Hashing

proach for typical values of M greater than 15. The number of generated pairs for both the Shape and Audio datasets does not exceed 10^9 pairs and therefore, can be processed in centralized settings. Conversely, the number of generated pairs for the Flickr dataset for intermediate values of M is high and cannot be handled in centralized settings. We recall here that the cost of storing one single pair is 10 bytes (2 integers for feature identifiers and 1 short for the collision frequency). The cost of processing 10^9 pairs is about 9.5 Gbytes. As a consequence, the default value of M for the Flickr dataset is fixed to 50. In the following, unless otherwise stated, the default value of M is 10.

In the following section, we use the Shape and Audio datasets to compare against the state-of-the-art technique applied in centralized settings. Results on the Flickr datasets are reported in 5.3 to evaluate the ability of our method to scale up in both dimensionality and dataset sizes.

5.2 Experiments in centralized settings

In this section, we first evaluate the overall performance of our method in centralized settings on only the Audio and Shape datasets (Section 5.2) to allow a fair comparison with the NN-*Descent* algorithm (Section 5.2) that could not run on the Flickr dataset.



(a) Shape (b) Audio

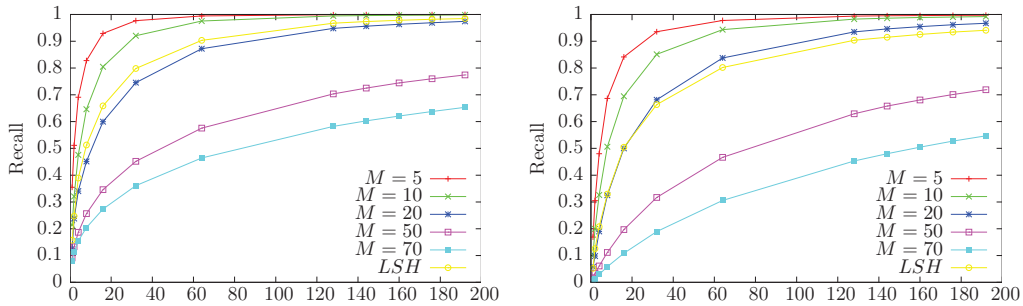
(c) Flickr

FIGURE 4.4 – Total number of collisions

Overall performance of our method and parameter discussion

Figure 4.5 summarizes the recall of our method on both Shape and Audio datasets for different hash functions and index sizes (i.e. the number L of hash tables used). The best results are observed for small values of M. The results also show high recall values even with a small number of hash tables ($L = 16$ and $L = 20$ respectively for 90% recall) whereas higher recall values require a higher number of hash tables (128 hash tables for $M = 10$, whereas only 64 hash tables are required for $M = 10$ for 99% recall).

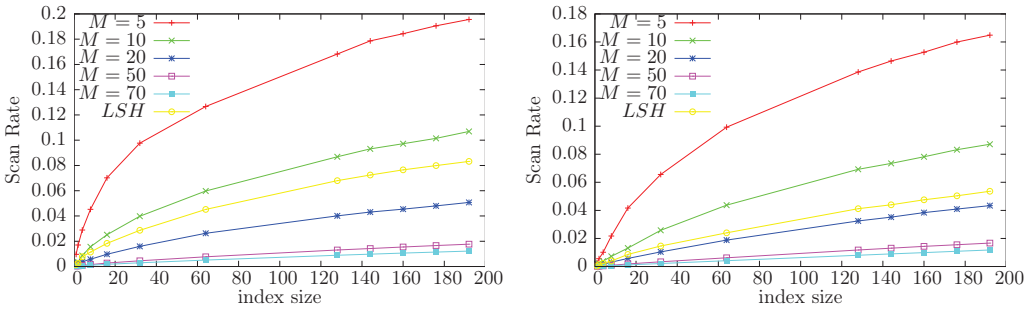
Note that high recall values can be achieved using different values of M. As discussed in 5.1, the higher M is, the fewer collisions are generated and the more hash tables are needed. Figure 4.6 plots the scan rate for different hash functions



(a) Shape (b) Audio

FIGURE 4.5 – Recall vs #number of hash tables used

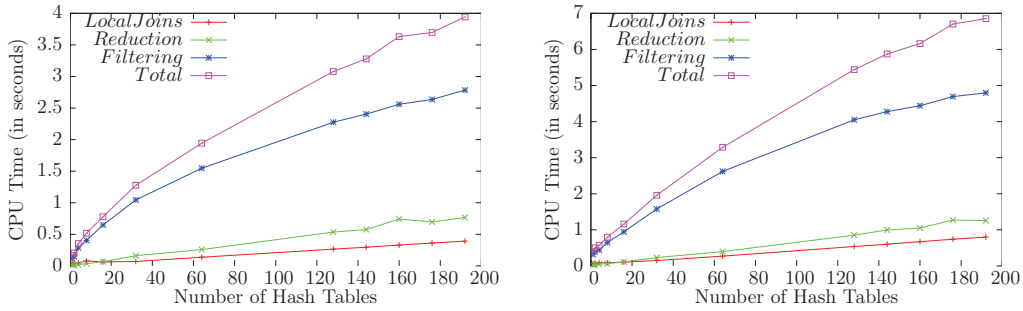
and index sizes. The results show that 99% recall can be achieved while considering less than 0.05 of the total $N * (N-1)/2$ comparisons (for $M = 20$) for both datasets. It is worth noticing that even with small values of M, and therefore low generalization properties ($M = 5$), the scan rate did not exceed several percent of the total number of comparisons. This suggests that intermediate values of M generate more accurate approximations of the KNN Graph as they require fewer comparisons for the same degree of accuracy. In practice, intermediate values of M with a high number of tables appears to be a reasonable trade-off between accuracy and approximation cost. Conversely, very high values of M degrade the accuracy of the KNNG approximation.



(a) Shape (b) Audio

FIGURE 4.6 – Scan rate variation vs #number of hash tables used

Figure 4.7 plots the total CPU time for different values of L ($M = 10$). For a better understanding of processing costs, we also report the running time of the different phases. The results show a correlation between the different phases. The greater the number of hash tables, the more collisions are generated along with irrelevant pairs.



(a) Shape (b) Audio

FIGURE 4.7 – Running Time - RMMH

In Table 4.4, we compare RMMH against LSH in both efficiency and effectiveness. RMMH clearly outperforms LSH with fewer hash tables. Actually, none of the LSH-based settings achieved better results than RMMH.

		Shape		Audio	
	L	Recall	Time	Recall	Time
RMMH	128	0.994	3.078	0.982	5.440
LSH	128	0.967	2.616	0.903	3.912
LSH	144	0.973	2.715	0.914	4.234
LSH	160	0.978	2.972	0.925	4.518
LSH	176	0.982	3.042	0.934	4.942
LSH	192	0.984	3.358	0.941	5.382

TABLE 4.4 – Total Running Time - LSH vs RMMH ($M = 10$)

Table 4.5 summarizes the Recall, Scan Rate and CPU time for different values of the minimum collision frequency considered in the approximate graph. Note that threshold values cannot exceed the number of hash tables used. The results show that low collision thresholds saved more than 50% of feature comparisons on

both datasets for less than a 4% recall loss.

	Shape			Audio		
thresold	Rec.	SR	CPU	Rec	SR	CPU
1	0,995	0,087	3,078	0,983	0,069	5,440
2	0,981	0,043	1,817	0,951	0,026	3,051
4	0,929	0,022	1,271	0,861	0,009	2,008
8	0,772	0,011	1,124	0,661	0,003	1,814

TABLE 4.5 – Impact of the filtering parameter ($M = 10$, $L = 128$)

Comparison with the State-of-the-art

The same datasets were used in [32] to evaluate the NN-*Descent* approach against Recursive Lanczos Bisection and LSH. We used the same NN-*Descent* settings in our set of experiments (($\rho = 0.5$) for speed and ($\rho = 1$) for accuracy).

We use $M = 10$ and two different thresholds for the post-processing phase ($t = 1$ (default) and $t = 2$) and up to 192 hash tables for high recall rates, as discussed in Section 5.2. Table 4.6 summarizes the recall and CPU time of both methods under those defined settings.

The results show that our approach yields similar results to the NN-*Descent* algorithm for both recall rates and CPU costs using the default threshold (i.e. a minimum collision frequency equal to 1). Although our method performs many fewer comparisons than the NN-*Descent* approach, the results show similar running times considering the Local Joins and Reduction costs. Higher threshold values are likely to further reduce the scan rate and CPU costs accordingly.

By putting a threshold on the frequency collisions, our method achieves both higher recall and faster speed ($t = 2$). Actually, our frequency-based approximation beats the NN-*Descent* high-accuracy setting in all cases. Here again, the results suggest that higher threshold values achieve better approximations of the KNN Graph.

As a conclusion, our method achieves similar or better performances than the most efficient state-of-the-art approximate K-NNG construction method in

centralized architectures. And in contrast to this method, our method has the advantage of being easily distributable and therefore much more scalable, as shown in the next section.

			Shape			Audio		
	ρ		Rec.	CPU	S.R.	Rec.	CPU	S.R.
NND	1		0,978	2,044	0,096	0,942	3,387	0,054
NND	0.5		0,958	2,33	0,057	0,903	4,834	0,033
	t	L	Rec.	CPU	S.R.	Rec.	CPU	S.R.
Ours	1	64	0,976	1,943	0,060	0,943	3,288	0,044
Ours	1	128	0,995	3,078	0,087	0,983	5,440	0,069
Ours	1	144	0,996	3,278	0,093	0,986	5,874	0,073
Ours	1	160	0,997	3,630	0,097	0,989	6,162	0,078
Ours	1	176	0,998	3,696	0,102	0,991	6,703	0,083
Ours	1	192	0,998	3,943	0,107	0,992	6,854	0,087
Ours	2	64	0,925	1,026	0,027	0,857	1,591	0,013
Ours	2	128	0,981	1,817	0,043	0,951	3,051	0,026
Ours	2	144	0,986	2,054	0,047	0,959	3,225	0,028
Ours	2	160	0,989	2,173	0,050	0,966	3,683	0,031
Ours	2	176	0,991	2,351	0,053	0,972	4,116	0,034
Ours	2	192	0,993	2,597	0,057	0,976	4,335	0,036

TABLE 4.6 – Comparison with State-of-the-art

Table 4.7 shows the impact of varying K (i.e. the number of Nearest Neighbors considered in the exact K-NNG) on both datasets ($L = 128$ and $M = 10$). It shows that high recall values can be obtained on both smaller($K = 1$) and larger graphs ($K = 20$) whereas a sufficiently large K is needed for the NN-*Descent* to achieve recall rates ($> 90\%$) as stated in [31].

K	1	5	10	20	100	Scan-rate
Shape	0.999	0.998	0.997	0.994	0.978	0.086
Audio	0.996	0.991	0.988	0.982	0.957	0.069

TABLE 4.7 – Recall for varying values of K

5.3 Performance evaluation in distributed settings

We recall that the experiments described here were carried out on the Flickr dataset (See Section 4.1).

Table 4.8, shows the impact of the *split local join* strategy on the number of *map* tasks for different values of parameter M. Despite the bucket balancing

achieved, the average maximum bucket size is still high. When the initial balancing of the hash tables is weak, large buckets are split into small balanced ones that fit the computational constraints resulting in a higher number of map tasks. On the other hand, the number of additional map tasks decreases as M increases to finally generate as many map tasks as the basic local join (*i.e.* the load balancing achieved by RMMH is already near perfect).

M	10	40	70	100
Gini	0.87	0.71	0.63	0.56
Basic Join	1161	1229	1260	1279
Split Join	9310	1256	1261	1280

TABLE 4.8 – Number of *map* tasks

In Table 4.9, we report statistics on *map* running times. Given a sufficient number of nodes, (i.e. $map\ slots \geq the\ number\ of\ map\ tasks$), the total processing time for a K-NNG construction would be of the same order of magnitude as the processing time of its longest map task.

M	10	40	70	100
Avg.	43	21	10	7
Worst	146	40	20	9

TABLE 4.9 – Map running time (in seconds)

Figure 4.8, displays the Recall/Precision curves for varying values of parameter M. Once again, the best results are observed for intermediate values of M between 15 and 70. This confirms the observations in [55] about the stability of this parameter and its expected optimal values. Very high values of M are likely to scatter similar objects and therefore, impact the recall. Conversely, low balanced buckets would lead to low precision rates. In the following, we fix the training parameter of RMMH to $M = 50$. In Figure 4.9, we report the Recall/Precision curves for varying numbers of hash tables. The results show that precision and recall rates increase along with the number of hash tables. In this experiment, we do not use any refinement step after constructing our approximate K-NNG, so relatively low precision could be drastically improved by a brute-force post-

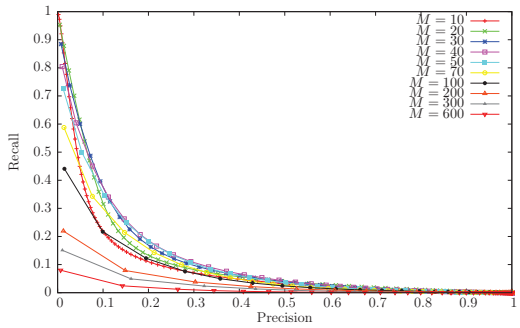


FIGURE 4.8 – ROC curve corresponding to the recall-precision curve on 128 tables

processing. To better evaluate the filtering capacity of our method, Figure 4.10

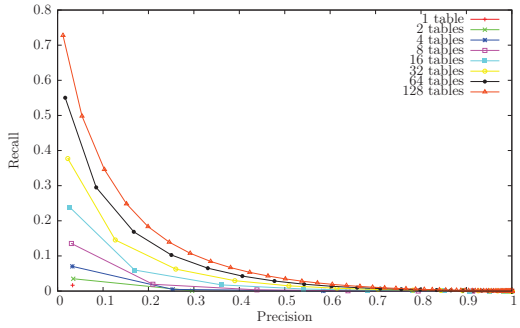


FIGURE 4.9 – ROC curve on Flickr dataset ($M = 50$)

plots the Recall/Scan-rate curves for increasing L values. Once again, it shows that increasing the number of tables always improve the trade-off between recall and scan-rate. With $L = 128$, 72% recall can be achieved while considering only 0.014 of the all pairs graph. Higher recall values could be achieved with more tables. But this is not required in many applications where approximate Nearest Neighbors can be as good as exact Nearest Neighbors from the document's content point of view [55].

Table 4.10 shows the impact of varying K on the Flickr dataset ($L = 128$ and $M = 50$). It shows that high recall values can be obtained on smaller graphs (e.g. K=1) and that very large graphs (K=1000) can still be well approximated with fair recall values about 50%.

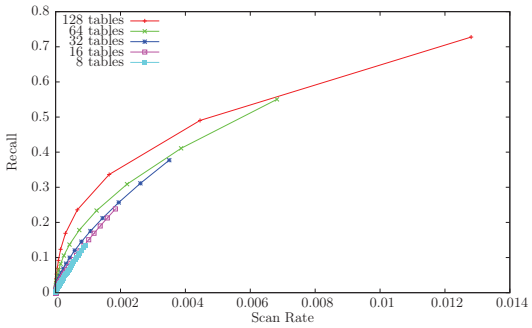


FIGURE 4.10 – Recall vs Scan-rate on Flickr dataset ($M = 50$)

K	1	5	10	40	100	$1K$	Scan-rate
Flickr	0.93	0.88	0.86	0.79	0.72	0.51	0.012

TABLE 4.10 – Recall for varying values of K

6 Conclusion

This chapter introduced a new hashing-based K-NNG approximation technique, which is easily distributable and scalable to very large datasets. To the best of our knowledge, no other work has reported full K-NN graphs results on such large datasets. Our study provides some evidence that balancing issues explain the low performances obtained with a classical LSH-based approach for approximating K-NN graphs. It also shows that using alternative new hash functions that handle hash table uniformity can definitely change those conclusions. Finally, we described a distributable implementation of our method under a MapReduce framework and improved the load balancing of this scheme through a split local join strategy avoiding memory overlaps.

In the following chapter, we show how the presented framework can be used in the context of event-based content suggestion through the construction of image similarity graphs. More importantly, we show how the collision scheme used can be leveraged to combine both the k-Nearest Neighbors Graph construction and the filtering step needed for the event k-Nearest Neighbors Graph construction.

Chapitre 5

Event-based Content Suggestion and Summarization in Social Media

Social media sites such as Flickr or Facebook contain large amounts of social multimedia documents relating real-world events. While some of their content might be interesting and useful, a considerable amount might be of little value to people interested in learning about the event itself. Applications such as event summarization, browsing and content suggestion would benefit from such identified content, ultimately improving the user experience. Selecting the most salient social media content for a particular event, however, is a challenging task, considering the fact that such User Generated Content is often distributed between a large number of different users. In this chapter, we address the general issue of selecting high quality content for an event.

We first present a new collaborative content based filtering technique for selecting quality documents for a given event (Section 1.1). We then extend our technique to support the more specific problems of event summarization (Section 1.2) and content suggestion (Section 1.3) in social media. Section 2, introduces a

scalable framework for building the event graph that we considered in this chapter. Section 3.1 reports results on the LastFM dataset used in the previous chapters.

1 Content suggestion and summarization in UGC

Recent studies have addressed the general issue of selecting relevant content, or summarizing an event. Very often, selecting the most interesting images involves some decision-making, based on various criteria. Alternatively, the problem of selecting relevant content can be reduced to an optimization problem under quality constraints [6, 93]. Nevertheless, these constraints vary greatly with the summarization context. Generally, state-of-the-art content selection and summarization techniques exploit the metadata associated to media such as time, location, title and description. In practice, such information is not always available or might be noisy.

To address the limitations of existing approaches, we leverage the social context provided by the social media to objectively detect moments of interest in social events. Our work is based on the assumption that, should a sufficient number of users take a large number of shots at a particular moment, then we might consider this to be an objective evaluation of interest at that moment. Of course, in such scenarios, location and time information provided with the contents have a major role to play. In practice, however, location and time information are not always available or might be noisy. In this chapter, we make use of the visual based event matching technique presented in Chapter 3, to fill in for bad or missing metadata associated to media.

1.1 Content Selection

For this content selection problem, we assume that we are given an event and a corresponding set of social media documents that are associated with the event, organized into records. Such identified record clusters can be obtained from the event k-NN Graph (Figure 5.1) described in Section 2 or by using metadata

associated to the media such as time, location and tags when available.

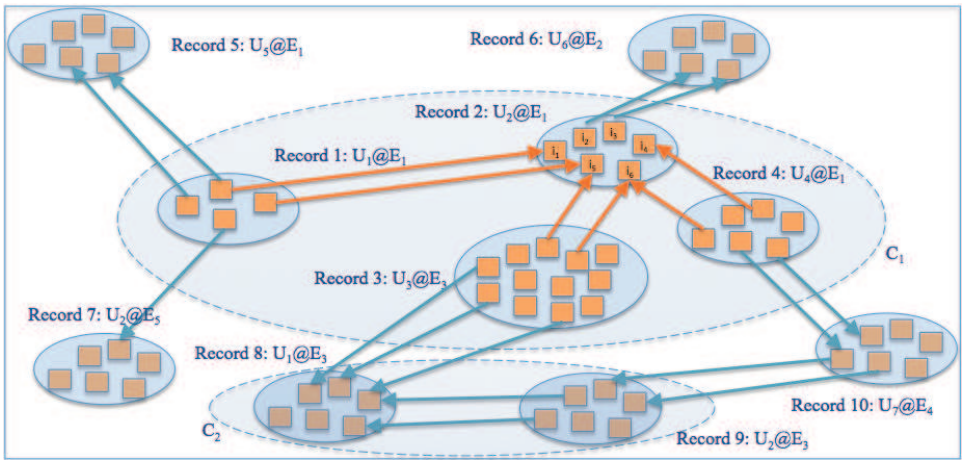


FIGURE 5.1 – A k-NN record graph of 10 event records.
C_1 = an event cluster of 4 event records related mostly to the social event E_1 ($U_1@E_1$, $U_2@E_1$, $U_4@E_1$ and $U_3@E_3$). C_2 = an event cluster of 2 event records related to the social event E_3 ($U_1@E_3$ and $U_2@E_3$).

Event clusters may, however, be noisy and contain records associated with some other events. This is particularly true in the case of co-located events where a set of people may be interested in the same event but also share images of other local events. Hence, records from different events are likely to share a subset of visually similar images and thus, appear within the same cluster. Figure 5.1 illustrates the situation described above. The event cluster C_1 contains 4 records related mostly to the social event E_1, it also includes an occurrence (Record 3) of the social event E_3. Conversely, distinct records may reflect different aspects of the event and, thus, be scattered between clusters. In our work, we assume that each record can be associated at most, with one single social event. However, our approach can generally be extended to handle less structured content.

Moreover, a single user record may relate to different events and thus, be associated with various events. For instance, the “my ICMR 2012” event record (Figure 5.2), containing images of my trip in the context of the International Conference on Multimedia Retrieval in 2012, may also include images of the “2012 Tiananmen celebration in Hong Kong” and images of our "trip to Lantau Island"

as well. Although these images relate to different events, they were taken in the context of "my trip to the ICMR conference" and may be considered as part of the same event. In this connection, images of the Tiananmen celebration might be of little interest to people who are seeking information about the ICMR 2012 event. Similarly, the ICMR pictures are likely to be of little value to people who are seeking information about the Tiananmen celebration ceremony.

FIGURE 5.2 – A photo collage of my 2012 ICMR photo album of co-located events.

Pictures of the ICMR banquet, however, are likely to interest people seeking information about the event. We define the *content selection problem* as follows :

Definition 3 *Given an event e and a set of associated social media documents D_e, our goal is to select a subset of documents from D_e that are the most relevant to the event, and which include useful details for people who are seeking information about the event.*

Our approach relies on the observation that widely covered moments are likely to reflect key aspects of the event as they reflect a common interest. Should a sufficient number of users take a large number of shots at a particular moment, then we might consider this to be an objective evaluation of interest at that moment.

Given a cluster C of n identified event records and their associated set of media documents I_c, our method counts, for each image $I \in I_c$, the number $S(I)$ of

temporally consistent visual matches with another image within a different record of the same cluster (*i.e.* the number of times that I contributed to a link with a record within the cluster). More formally, let G be the graph having elements from I_c and whose edges link pairs of temporally and visually consistent images from I_c. The $S(I)$ score represents the in-degree centrality of I. The result is a ranked subset of images of I_c.

1.2 Event Summarization

We formalize our event summarization problem as that of producing a ranking on the event media documents. Specifically, given a cluster C_i (relating to an event E_i) of n event records $U_j@E_i$ and their associated set of media documents I_c, we first compute the $S(i)$ score for each element in I_c. We then select the top-K documents accordingly as a candidate set for generating an event summary. In our experiments, since the average number of images per event cluster is relatively small (from tens to hundreds), we make K big enough to include all the images in the event clusters. Going back to the record map illustrated in Figure 5.1, the E_1 summary is generated by ranking the images of C_1 records in decreasing order of their in-degree centrality.

Alternatively, the resulting set may be post-processed to produce customized event summaries, thus, improving the user experience. In practice, we provide users with a set of predefined filters, so that, for example, the removal of visually similar images (possible using the images K-Nearest Neighbors Graph) or maximize the time span (when temporal information is available) of the images for a wider coverage of the event. Since the number of images retained may decrease, we refer to the size of the pruned summary as $S_{summary}$.

1.3 Content Suggestion

Here, the goal is to present a given user only documents that provide additional information about the event. Given a set of N_q images (i.e. a record of a user), the

recommendation system first identifies the corresponding event and then, returns a ranked list of images from the repository.

In practice, the event record is submitted to the system and matched with records from the repository using the visual-based event matching technique described in Section 2. To suggest images to the user, we follow two possible scenarios.

In one scenario, we do not have any information about the retrieved records. In this scenario, the set of suggested images is that of the first retrieved event record. Depending on the system requirements, a threshold on the retrieved records can be tuned to improve the overall precision, respectively recall, of the recommendation system.

Alternatively, all event records in the repository are clustered in an offline phase. Finally, the recommendation system, returns the list of images of the identified event records ordered by their score (Section 1.1).

To illustrate both scenarios, let us consider the record graph in Figure 5.1 and a query record that matches with a record from C_1. In the first scenario, the set of suggested images is that of the matched record. In the second scenario, the set of suggested images is expanded to include images from C_1 ordered by their decreasing score.

In both scenarios, images which are visually similar to the query images are removed from the answer set, as they would not provide any additional information.

2 Building the Records Graph

The experiments in Section 4.2 show that big values of k are needed to achieve both good precision and recall. However, a large proportion of the records retrieved are discarded while applying the spatio-temporal constraints the registered records. Here, the idea is to discard records with large spatial and/or temporal offset from the query record prior to the geo-temporal re-ranking step (Step 3) hence, combining both the visual matching (Step 1) and prior filtering (Step 4)

steps.

In Chapter 4, we presented a framework for large scale nearest neighbors graph construction. Candidate pairs of visually similar images are first produced using our hash-based Nearest Neighbors collision scheme (Step 2). A threshold of the number of collisions is then used in order to keep only the most similar pairs of images (Step 4). Eventually, a refinement step on the the remaining pairs is performed to determine the k-nearest neighbors of each image using a distance-based similarity function.

Although the number of generated pairs is relatively low, it is still has an impact on the amount of data transferred. The idea here is to limit the set of emitted pairs to those that fit the specified spatio/temporal constraints (Step 4) and thereby also limiting the number of similarity computations performed in the refinement step. Algorithm 3 gives the pseudo-code for the enhanced local-join.

Algorithm 3 *Visual and temporal Local Join*

Require: bucket $b = \{id_i\}_{1<i<n_b}, start, end.$
Ensure: local collision set C.

```
1: C ← ∅
2: for i ← start, ..., end do
3:     for j ← i + 1 ..., n_b do
4:         if P^{b_i} − P^{b_j} ≤ δ_max then
5:             C ← C ∪ (b_i, b_j)
6:         end if
7:     end for
8: end for
```

The approximate event Nearest Neighbors Graph is computed by applying Step 2 and 3 on the approximate image k-Nearest Neighbors obtained using the modified local join algorithm.

3 Experiments

We evaluated our content selection technique on an 828,902 *Flickr* images dataset, the same as was used in [94]. We first describe the experimental settings

(Section 3.1). Experimental results are reported in Section 3.2.

3.1 Experimental setup

Data

The KNN-Graph on the records set was built using the exact image's Nearest Neighbors on the global features and the default δ and θ parameters as in Section 4.2.

From the full set of 34,034 events, the only events we kept were those having at least 2 related records in the dataset. The resulting graph contains 11,785 event records from 4,525 different sub-events. We used the LastFM tags associated to the images to build event records clusters and compute the $S(i)$ score of the related images.

Evaluation

We conducted a user-centric evaluation on 10 different subjects. Each user was asked to evaluate a set of 20 event summaries chosen at random from a set of 168 events, each having at least 5 associated event records. A 1 to 5 scale was used to score the overall quality of the summary, where a score of 5 signifies strong relevance and clear usefulness, and a score of 1 signifies no relevance and no usefulness. Similarly, a 1 to 5 score was used to score the images of the summary individually. The number of images displayed was limited to the top 7 ranked images so that summaries could fit into web browsers. For each event, we report the event summary and the average of the recommended images. Figure 5.3 illustrates the web-based application used for this purpose.

The KNN-Graph on the records is evaluated using the optimal values from sections 4.2 and 5 (i.e. δ_{max}=86.400, K=3000, θ=1800, $M = 50$) as described in Section 4.1.

FIGURE 5.3 – Snapshot of the user-centric evaluation GUI

3.2 Results

We first evaluate the ability of the proposed method to suggest relevant content. Figure 5.4 shows the score distribution of the suggested images. The results show that 39% of the suggested images were rated with the highest score while only 5% had the lowest. Overall, 68% of the scored images were judged good enough to represent the event they belonged to.

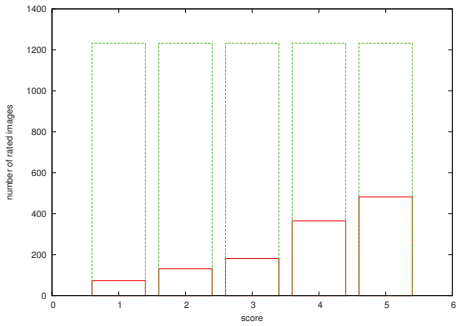


FIGURE 5.4 – Score distribution of the suggested images

Looking at the results in more detail, we concluded that, without much surprise, the worst rated images are generally those displaying only a few people not directly participating in the main event itself (friends of the photographer, a

lunch break, etc.) or images of very poor quality. On the other hand, the top-rated images are usually good quality images where the artist(s) is(are) clearly visible and/or where the scene presents a specific interest. This can be observed in Figures 5.5 and 5.6, where close-up photos of the artists rated considerably higher than photos of the scene and the venue as well as these 3 events were co-located in Hong-Kong at the same period. Although the Pukkelpop festival art work shot, captured dining the event (Figure 5.5 first image with past events dates), seems to have no particular interest at first glance, it rated 3,33 on average as it provides information about the past events.

FIGURE 5.5 – Pukkelpop Festival 2007 summary. The first image was rated at 3.33 on average whereas the remaining images rated at 4.33, 4.33, 4 and 4.33 on average, respectively.

In Figure 5.7, we compare the event summary score (given by users) to the image-based event score (the average score of the suggested images). The results show that the two scoring methods yield very similar results. None of the rated events had a score of 1, while 65% to 73 % scored higher than 4. Although the two reported scores were similar, the results show some variation. A higher image-based score, for instance, reflects a limited event coverage despite the quality of the suggested images (Figure 5.8). Conversely, a higher summary score reflects wide coverage of the event (Figure 5.6). Still, the average summary score was 3.75, respectively 3.94, which reflects the effectiveness of our scoring approach and

FIGURE 5.6 – Haldern Pop Festival - August 13-19, 2009 Summary. All of the images were rated at 4.5 on average.

subsequently, the summarization technique.

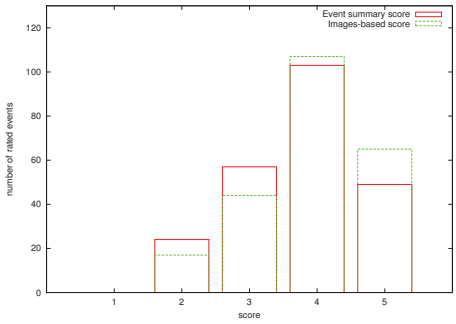


FIGURE 5.7 – Event summary vs image-based score distribution.

Detailed user statistics of the user study are presented in Table 5.1. The results show consistency between the images and the events score for each user. Although the events scores ranged from 2 to 5 for most users, the average event score was high. This suggests that our approach was able to present comprehensive summaries for users, who may well have different requirements and expectations.

In Figure 5.9 we report the average scores for varying event cluster sizes. The results show that the effectiveness of our technique increases with the size of event

FIGURE 5.8 – Radiohead @ Victoria Park - June 24, 2008 Summary. The event summary was rated at 3 while the image based score was at 2.

	Avg Image score	Avg Event score	Worst Event rating	Best Event rating
U_1	4.06	4	2	5
U_2	3.69	3.63	2	5
U_3	3.97	3.85	3	5
U_4	3.73	3.7	2	5
U_5	4.38	4	2	5
U_6	4.09	4	3	5
U_7	3.66	3.72	2	5
U_8	3.67	3.8	2	5
U_9	3.34	3.31	2	5
U_{10}	4.09	3.95	2	5
Mean	3.87	3.79	2.2	5

TABLE 5.1 – User-centric evaluation of the image relevance scores

clusters (*i.e.* the number of records) as more information is available. The results also show that only a relatively small number of records is needed to generate a representative summary of the event.

Figures 5.10 and 5.11 show the impact of the near duplicate pictures removal step (Section 1.2). Duplicate images are removed within the summary and possibly replaced by the next images in the ranked set of the selected images, thus providing a wider and enhanced coverage of the event.

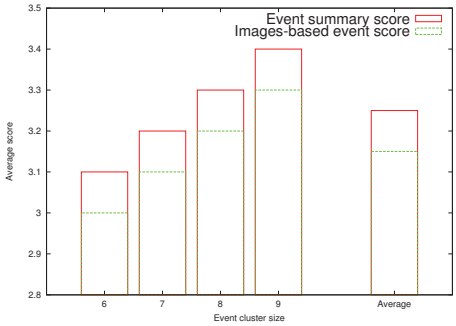


FIGURE 5.9 – Average score per event cluster size

FIGURE 5.10 – An event Summary without duplicate pictures removal filter

FIGURE 5.11 – An event summary showing the impact of the duplicate pictures removal filter

Building the Events K-NN Graph

In Figure 5.12, we report the mean average precision for varying values of k and the number of hash functions used. The results show that MAP values increases along with the number of hash functions used, as the number of generated collisions increases. Similarly, the results show that the mean average precision increases rapidly along with the number of Nearest Neighbors retrieved, to level out for high values of k. In our experiments, the number of Nearest Neighbors retrieved is relatively low resulting in a constant mean average precision for high values of

k. In the following, unless otherwise stated, the number of hash tables used is 10.

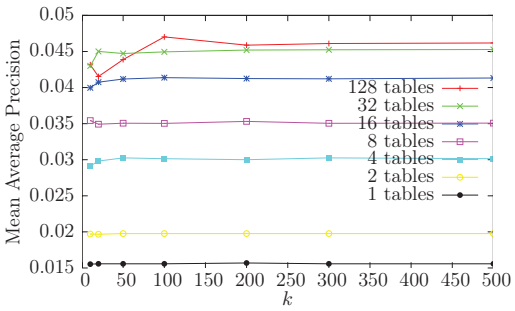


FIGURE 5.12 – Mean Average Precision vs k

In Figure 5.13, we report the precision and recall curves for increasing values of k. The results show that the recall increases rapidly along with the number of Nearest Neighbors retrieved. Converserly, the overall precision decreases as the number of retrieved neighbors increases. Overall, the reported results show a large gap between precision and recall. Most importantly, Figure 5.13 show that the recall increases twice as fast as the precision decreases. This suggests that higher recall can be achieved without a significant loss in precision. In the following, we study the impact of the selectivity of the hash functions used on both precision and recall.

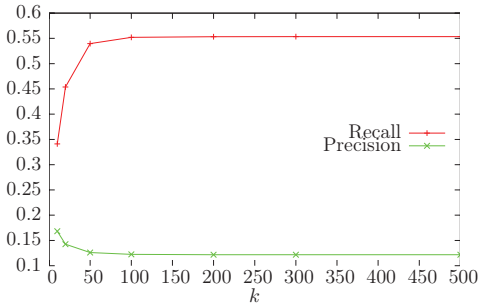
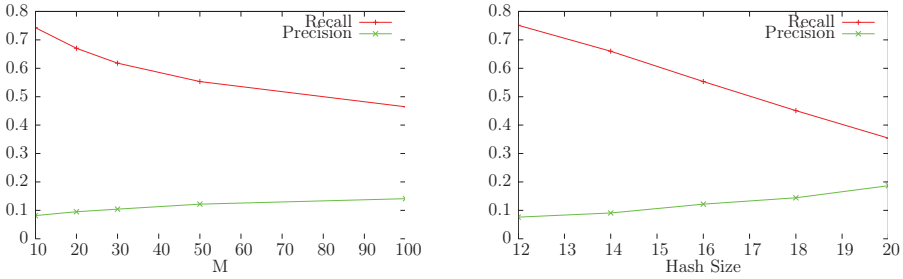


FIGURE 5.13 – Recall and Precision vs k

Figure 5.14 shows the influence of the selectivity of the hash functions used on the recall and precision. The results show that high recall values can be achieved without a substantial loss in precision. Specifically, Figure 5.14a shows higher recall

for decreasing values of M as as the selectivity of the hash functions decreases. Similarly, Figure 5.15 shows higher recall for decreasing size of the hash functions used as multiple buckets are merged.



(a) Recall and Precision vs M parameter

(b) Recall and Precision vs Hash Size ($M = 50$)

FIGURE 5.14 – Influence of the hash functions selectivity on the recall and precision

Figure 5.15, shows the combined effect on both recall and precision. Although recall improved significantly, the achieved precision is sill low. In the following, we fix the training parameter of RMMH to M = 10 and the Hash Size of the hash functions used to 12 and study the impact of the filtering parameter on both recall and precision.

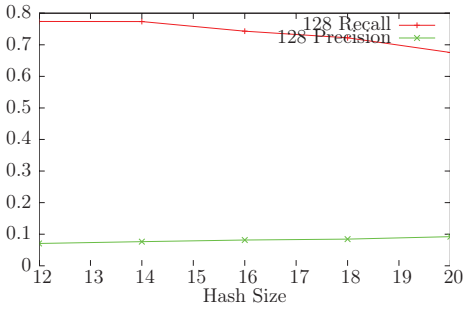
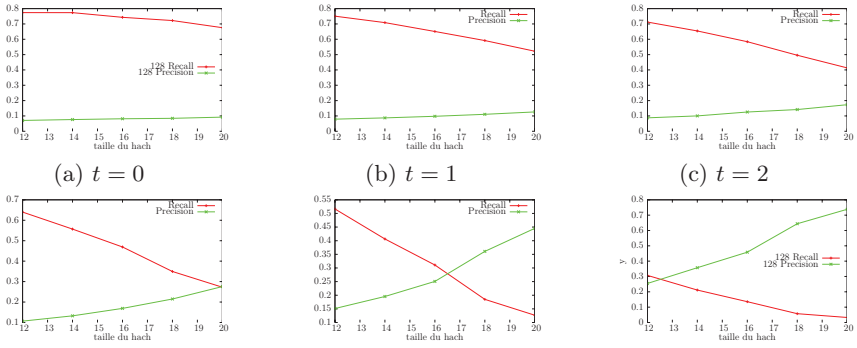


FIGURE 5.15 – Recall and Precision vs Hash Size ($M = 10$)

Figure 5.16 shows the impact of the filtering parameter on both recall and precision. By putting a threshold on the frequency collisions, our method achieves higher precisions but at lower recall resulting in a better balance between precision and recall.



(a) $t = 0$ (b) $t = 1$ (c) $t = 2$

(d) $t = 4$ (e) $t = 8$ (f) $t = 16$

FIGURE 5.16 – ROC curve for various collisions thresholds

Such performances are clearly acceptable from an application point of view as shown in Figure 5.2. Overall, in the best case, our method is able to identify the correct event with a 56% success rate and to suggest the correct event tag over 5 suggestions with a 92% success rate.

TABLE 5.2 – Suggestion rates

# of suggested events tag	1	2	3	4	5	6	7	8	9	10
Suggestion rate	0.56	0.75	0.84	0.88	0.89	0.9	0.91	0.9	0.91	0.92

4 Conclusion

Events in social media often have vast amounts of associated content. In order to avoid overwhelming their users with too much information, social media sites need to select and prioritize content.

In this chapter, we presented a new content-based filtering technique to select high quality content. Unlike state-of-the-art methods, our record-based technique provides an objective evaluation of moments of interest during social events. A user-centric evaluation revealed that some users tend to prefer stage photos while others see more value in more diversified content. Overall, the proposed technique

has performed well, reporting the most captured moments for a set of users. We argue that such information can be used to characterize communities of users, and more generally, social networks.

Chapitre 6

Related Work

This chapter reviews the literature that is relevant to this dissertation. Section 1 describes efforts on event identification in social media related to the event retrieval task presented in Chapter 3. Section 2 discusses related research on organizing and presenting social media content, including content suggestion and summarization, which we addressed in Chapter 5. Section 3 and Section 4 respectively, provide an overview of approximate nearest neighbors search techniques and large-scale k-Nearest Neighbors Graph construction that we considered for the event graph construction framework presented in Chapter 4.

1 Event Identification in Social Media

While earlier studies aiming at event discovery were based solely on text analysis and essentially focused on news documents [61, 63], more recent work on social media has been able to take advantage of richer multimedia content, while having to cope with the challenges that such a benefit entails.

In [4], the authors present a clustering framework to group images based firstly on geographical coordinates and then visual features to depict different views of the same scene. Similarly, in [77], a framework to detect landmarks and events to improve the user browsing and retrieval experience is proposed. The work presen-

ted in [36] attempts to identify public events by using both the spatio-temporal context and photo content. Although these methods performed well on real word datasets, their scope remains conceptually limited due to the properties of the EXIF data, which considerably restricts query formulation. Moreover, despite the fact that such properties are becoming more widespread, they are far from being universally available, notably in professional devices, or removed (such as for Flickr and Facebok).

Several recent studies have tried to compensate for such missing information. For instance, [11] presents a classifier-based method, where items that are geotagged are used to build a set of initial clusters that correspond to events. The items of each identified cluster are then used to train a classifier that augments each cluster with non-geotagged items. More recent work also extended item description with information from user-supplied tags along with external data sources and APIs such as the Google MAP API. In [82], the authors make the assumption that all items that have been uploaded by the same user in the same day belong to the same event. Such heuristics make sense. However, their use may introduce some bias in return.

Other related efforts have used some online sources to retrieve structure information that is related to an event. The work of [91] exploited the user context to annotate the images according to four event related facets (where, when, who, what) by a graph model that uses the Wordnet [71] ontology. In [99], a sequence of clustering and filtering operations is applied. The textual, temporal and location features are first used to cluster images. The resulting clusters are then filtered with regard to the temporal, location and textual constraints. A visual classifier is then used to filter clusters. This final step, however, requires manual labeling of images. Similarly, in [66], the authors propose an approach that builds a classifier using explicit event descriptions from online catalogues and performs post-processing on the visual features to clean the classified data. In [76], the clustering step is based solely on the location and temporal information. Each event produced by the clustering step is then enriched by making use of the metadata of the photos

associated with it, including pictures by the same users or within a fixed radius of the venue. In [67], the description of events and their associated multimedia items is retrieved from structured online sources and expressed using the LODE [90] and Media Ontology respectively. Here, multimedia documents that contain specific machine tags are used to train classifiers which are then used to prune results from general textual queries. Although ontologies provide a common description of real world-events, their practical use is still limited by the number of searchable properties as well as the lexical ambiguity of textual based queries. In [82], the authors introduce an approach to detect photos belonging to the same event which are distributed between several friends and albums in Facebook using visual, tag-based, friendship-based and structural-based features.

In general, many of the approaches that have been proposed to tackle the problem of event identification in multimedia collections have used some form of online source to retrieve structured information that is related to either the event or media. While this is acceptable if it leads to an enhancement of the results, it may not always be possible as most social events do not have a formal description in some online source. Therefore such methods should only be used for pruning results. Although some of these efforts make use of certain visual properties, very few rely primarily on visual features. Our work differs in that it essentially relies on visual features to identify event-related items while incorporating additional external information when such information is available.

2 Event summarization

In the computer vision community, [96] and [72] provide an extensive review of key-frame extraction and video summarization. In [56], broadcasted videos of a an entire basketball season in the USA and the corresponding metadata are used to create summary videos from different aspects, like summaries of the whole championship, of only one team or even a single player. In [33], the authors present an approach for summarizing rushes video based on the detection of repetitive

sequences, using a variant of the Smith-Waterman algorithm to find matching subsequences.

Other recent efforts have addressed the problem of presenting and summarizing web images. In [98], the authors create a “picture collage”, a bidimensional spatial arrangement of the input that maximizes the visibility of salient regions. Rother et al. [84] summarize a set of images with a “digital tapestry”. A large output of images is produced, stitching together salient and spatially compatible blocks from the input image set. In both cases, however, the set of images to appear has already been selected, and the visual layout is to be determined.

In social media, early work focused on extracting quality Twitter messages [25, 30] and summarizing or otherwise presenting Twitter event content [6], an effort related to ours but using fundamentally different data. In [25, 30], the authors analyzed Twitter messages corresponding to large-scale media events to improve event analytics and visualization. In [6], Becker et al. address the problem of selecting tweets with regard to quality, relevance and usefulness.

Selecting the most representative social media documents from large collections of social media documents is becoming a prominent issue in the multimedia community. Early work focused solely on metadata associated to media [21, 47, 88, 59]. In [29], the authors make use of community annotations, such as ratings and the number of views, to produce video summaries of social events using both video and image content.

In [21] a hierarchy of images is constructed using only textual caption data, and the concept of subsumption. Jaffe et al. [47] summarize a set of images using only tags and geotags. Here, the authors use the correlations between geotags and tags to produce "tag maps", where tags and related images are placed on a geographic map at a scale corresponding to the range over which the tag appears. In [88], the author, relies on Flickr tags, which are typically noisier and less informative than captions. All of these approaches could be used to further organize our summaries. However, none of them take advantage of the visual information in the images to fill in for bad or missing metadata. Hence, in [59], the authors propose a method

to generate representative views of landmarks by diversifying image features and user tags.

Content summarization, however, turns out to be a very subjective process. In [87], Savakis et al. show that selecting personal photos from a collection depends greatly on user-preferences. In [93], Sinha et al. address the problem of summarizing personal photos present in web archives or personal storages with high quality, diversity and coverage constraints. Here, the authors reduce the problem of selecting images from photo collections to an optimization problem under quality, diversity and coverage constraints. A framework, based on spatial patterns, for automatically selecting a summary set of photographs from a large collection of geo-referenced photos is presented in [47]. Here, the authors make the assumption that more photographs are taken at locations that provide views of some interesting object or landmark by a large number of photographers. Although these efforts make use of certain visual properties, very few make use of the social media information associated to media, such as tags of individuals [82] and ownership [94].

More generally, research on content selection, and event summarization has benefited from recent work on event identification and retrieval in social media. In a notable effort, Liu et al. present a method combining semantic inferencing and visual analysis to automatically find media (photos and videos) illustrating events. In [94] , we presented a new visual-based technique for retrieving events in photo collections, typically in the context of User Generated Content. Given a query event record, represented by a set of photos, the proposed method aims to retrieve other records of the same event, typically generated by distinct users. One advantage of this approach is that it essentially relies on visual features to match records of the same event while incorporating additional external information when such information is available.

3 Large-scale k-NN Graph construction

Initially, the K-NNG problem can be seen as a nearest neighbors search problem where each data point itself is issued as a query. The brute-force approach, consisting of N exhaustive scans of the whole dataset, has the cost $O(N^2)$. Its practical usage is therefore limited to very small datasets. Building an index and iteratively processing the N items in the dataset with approximate Nearest Neighbors search techniques is an alternative option that might be more efficient (Section 4).

In addition to the usual approximate Nearest Neighbors search methods, some recent studies focus more specifically on the K-NNG construction problem as a whole, i.e. not by processing iteratively and independently the N top-K queries, but by trying to exploit shared operations across all queries. In the text retrieval community, recent studies [5, 107] focused on the ϵ-NNG construction in which one is only interested in finding pairs whose similarity exceeds a predefined threshold. In [107], the authors present a permutation based approach both to filter candidate pairs and to estimate the similarity between vectors. However, their approach is only applicable on sparse vectors. Very recently, Dong et al. [31], proposed the NN-*Descent* algorithm, an approximate K-NNG construction method purely based on query expansion operations and applicable to any similarity measure. The algorithm starts by picking an approximation of K-NN for each object, it iteratively improves that approximation by comparing each object against its current neighbors' neighbors, and then stops when no improvement can be made. Their experiments show that their approach is more efficient than other state-of-the-art approaches. However, designing an efficient distributed version of this method is not trivial, limiting its practical scalability as it requires the entire dataset to be loaded into a centralized memory. A comparison with their results was provided in Chapter 4.

4 Nearest Neighbors search

Early tree-based indexing methods for Nearest Neighbors (NN) search such as R-tree [43], SR-tree [57], M-tree [20] or more recently cover-tree [9] return accurate results, but they are not time efficient for data with high dimensionality [100].

4.1 Curse of dimensionality

A particular but well-studied case of the nearest neighbor search problem is in the Euclidian space where the data lives in a d-dimensional space $\mathbb{R}^d$ under the Euclidean distance function.

When $d = 1$, predecessors queries can be used to efficiently perform nearest neighbour queries. A straightforward but efficient solution is to sort the data at indexing time, and then, perform a binary search at query time. This achieves linear spacey and polylograithmic time complexity.

The $d = 2$ case leads to one of the most classical structures in computational geometry, the *Voronoi diagram*[26]. Here, the plane is partitioned into polygonal regions, each representing the set of points that are closer to a point from the dataset to any other point from the dataset. At query time, one just need to locate the region containing a given query.

While the latter approach achieves $O(n)$ and $O(n \log(n))$ space and time complexity, respectively, its generalisation have $O(n^{\lceil d/2 \rceil})$ space complexity. In practice, such a space bound is impractical for datasets of a few million points for $d \geq 3$.

Several data structures have been proposed for low values of d. Kd-trees, introduced were first such structure in 1975 by Bentley et al. [8]. In [35], such a structure is used to accelerate k-nearest neighbour queries using ball-rectangle intersection tests. Ever since, many approximate NN methods were then proposed including randomized kd-trees [92], hierarchical k-means [73] or approximate spill-trees [65, 50]. Although these methods provided little improvement over a linear time algorithm that compares a query to each point from the database, they are

not time efficient for data with high dimensionality [100].

Since, several approaches have been proposed to overcome space and time limitations using approximation. One of the most popular approximate nearest neighbor is LSH. In that formulation, we are no longer interested in the exact k-nearest neighbors trading accuracy for time and space efficiency.

4.2 Approximate similarity search

Approximate nearest-neighbor algorithms have been shown to be an interesting way of dramatically improving the search speed, and are often a necessity [106, 19]

Locality-Sensitive Hashing

One of the most popular approximate nearest neighbor search algorithms used in multimedia applications is Locality-Sensitive Hashing (LSH) [37, 46]. The basic method uses a family of locality-sensitive hash functions composed of linear projections over randomly selected directions in the feature space. The principle is that nearby objects are hashed into the same hash bucket with a high probability, for at least one of the hash functions used. LSH has achieved very good time efficiency for high dimensional features and has been successfully applied in several multimedia applications including visual local features indexing [58], songs intersection [14] or 3D object indexing [70]. Following this success, hashing methods have been gaining increasing interest.

Multi-Probe LSH

One drawback of the basic scheme is that, in practice, it requires a large number of hash tables (L) to achieve good search accuracy. In [74], Panigrahy et al. proposed an entropy-based LSH scheme to reduce the number of hash tables required by using both the original query point and randomly perturbed nearby points as additional queries.

To make better use of a smaller number of hash tables, Lv, et al. [68] not only considers the bucket pointed by the query point, but also examines "nearby" buckets. Here, instead of perturbed query objects, the authors generates perturbed hash tables. However, this method still suffers from the need of building hash tables at different radiuses in order to achieve good search accuracy.

Whereas the latter are based on the simple likelihood criterion that a given bucket contains query results, in [54], the authors define a more reliable a posteriori probabilistic model taking account some prior about the queries and the searched objects. This prior knowledge allows a more accurate selection of buckets to be probed.

So far, hashing techniques are categorised into two groups : **Data independent hashing functions** in which the hashing function family is defined uniquely and independently from the data to be processed [18, 83, 89, 48] and more recently in [52] and [51], and **data dependent hashing functions** in which the hash functions rely on some features sampled in the dataset [101, 60, 78, 86, 55]. Efficiency improvements of data dependent methods over independent ones have been shown in several studies [49, 101, 86, 55]. RMMH [55], the family used in this work, was designed to overcome two limitations of previous data dependent methods : (i) it is usable for any Mercer Kernel (ii) it produces more independent hashing functions.

Whereas most of the latter approaches have tackled the approximate nearest neighbours problem in the euclidian space some recent work addressed the problem using x^2 distance which is believed to achieve better results in image retrieval context. In [41, 40], the authors present a new LSH scheme adapted to x^2 distance for approximate nearest neighbours search in high-dimensional spaces that achieves better accuracy than euclidean scheme for an equivalent speed, or

equivalent accuracy but with a high gain in terms of processing speed.

Recently, there have been several efforts to improve the load balancing of the generated hash functions. For unsupervised hashing, principled linear projections like PCA Hashing (PCAH) [97] and its rotational variant [39] were suggested for better quantization rather than random projections. Nevertheless, only a few orthogonal projections are good for quantization as the variances of data usually decay rapidly, as pointed out by [97]. In [45], the authors present a novel hypersphere-based hashing function, spherical hashing, to map more spatially coherent data points into a binary code compared to hyperplane-based hashing functions. Intuitively, hyperspheres provide much stronger power in defining a tighter closed region in the original data space than hyperplanes. For example, while d + 1 hyperplanes are needed to define a closed region for a d-dimensional space, a single hypersphere can form such a closed region even in an arbitrarily high dimensional space.

Conclusion

Summary of Contributions

As people continue to author and share event-related content in social media, the opportunity for leveraging such information increases. Social media web sites such as Flickr and Facebook, provide a playground not only for people to publish their content but also for applications that build on these useful sources of information. While some of event related content might be interesting and useful, a considerable amount might be of little value to people, ultimately impacting the user experience.

In this dissertation, we presented a visual-based event matching paradigm which serves as a stepping stone for various applications that build on events, and their associated documents, in social media. In Chapter 3, we addressed the problem of identifying events in social media web sites. By linking different occurrences of the same event, we can annotate the query with tags from previously identified and/or annotated occurrences. Ultimately, linking different occurrences of the same event would enable rich search and browsing of social media events content. Specifically, linking all the occurrences of the same event would provide a general overview and description of the event.

To avoid overwhelming applications, or users, with unmanageable volumes of event-related content, we presented a new collaborative content-based filtering

technique for selecting quality documents for a given event (Chapter 5). Subsequently, we addressed the more specific problems of event summarization and content suggestion in social media.

To improve our content selection framework, we developed a scalable and distributed framework for k-Nearest Neighbors Graph construction (Chapter 4) based on RMMH. Our work provides some evidence that balancing issues explain the low performances obtained with a classical LSH-based approach for approximating K-NN graphs. It also shows that using alternative new hash functions that handle hash tables uniformity can definitely change those conclusions. We finally described a distributable implementation of our method under a MapReduce framework and further improved the load balancing of this scheme through a split local join strategy to accommodate memory requirements.

Future Work

Identifying communities

By linking different occurrences of the same event, we can identify communities of users who share a comment interest in a specific event or a particular group of events, ultimately extending the event experience and allowing users to socialize and share their experience.

Collaborative event recommendation

Obviously, people attending the same event are likely to have similar tastes and preferences. By connecting users accordingly, we can discover more complex relationships between users, as well as the events themselves. A user graph could for instance be obtained straightforwardly from our event records graph.

A dedicated framework for k-NN Graph Construction

The k-NN graph construction framework developed in this dissertation (Section 3) is at the core of our content suggestion and event identification techniques.

Although the technique presented is scalable, the Hadoop-based implementation suffers from some technical limitations. Most importantly, the C++ API requires data to be serialized in order to be sent to and from the mappers respectively, the reducers. A dedicated framework would allow data to be handled natively, hence, improving the overall performance.

Further work remains to be carried out on automatic parameters tuning and varying hash sizes through a rigorous theoretical analysis of our method. Other perspectives include : query expansion strategies, hash functions evaluation and metadata management.

Finally, the technique presented could extend a large pool of existing graph and network analysis methods to large datasets without an explicit graph structure. Further work will be carried out towards extending our framework to support large-scale data mining techniques.

Distributed event records Graph construction

So far, we have presented a scalable framework for k-NN Graph Construction. However, the event records Graph construction, from the image similarity graph, is still centralized. A short-term perspective of this work is to distribute the construction of the event records Graph.

New record similarity metrics

So far, we have only considered the use of temporal meta-data, taking into account the fact that spatial information is rarely available. Further work should be carried out to include additional, more abundant meta-data such as textual annotations.

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Découverte d'événements par contenu visuel dans les médias sociaux

Mohamed Riadh TRAD

RESUME : L'évolution du web, de ce qui était typiquement connu comme un moyen de communication à sens unique en mode conversationnel, a radicalement changé notre manière à traiter l'information. Des sites de médias sociaux tels que Flickr et Facebook, offrent des espaces d'échange et de diffusion de l'information. Une information de plus en plus riche, mais aussi personnelle, et qui s'organise, le plus souvent, autour d'événements de la vie réelle. Ainsi, un événement peut être perçu comme un ensemble de vues personnelles et locales, capturées par différents utilisateurs. Identifier ces différentes instances permettrait, dès lors, de reconstituer une vue globale de l'événement. Plus particulièrrements, lier différentes instances d'un même événement profiterait à bon nombre d'applications tel que la recherche, la navigation ou encore le filtrage et la suggestion de contenus.

L'objectif principal de cette thèse est l'identification du contenu multimédia, associé à un événement dans de grandes collections d'images. Une première contribution est une méthode de recherche d'événements basée sur le contenu visuel. La deuxième contribution est une approche scalable et distribuée pour la construction de graphes des K plus proches voisins. La troisième contribution est une méthode collaborative pour la sélection de contenu pertinent. Plus particuliérement, nous nous intéresserons aux problèmes de génération automatique de résumés d'événements et suggestion de contenus dans les médias sociaux.

MOTS-CLEFS : Recherche d'événements, résumés d'événements, graphes des plus proches voisins.

ABSTRACT : The ease of publishing content on social media sites brings to the Web an ever increasing amount of user generated content captured during, and associated with, real life events. Social media documents shared by users often reflect their personal experience of the event. Hence, an event can be seen as a set of personal and local views, recorded by different users. These event records are likely to exhibit similar facets of the event but also specific aspects. By linking different records of the same event occurrence we can enable rich search and browsing of social media events content. Specifically, linking all the occurrences of the same event would provide a general overview of the event.

In this dissertation we present a content-based approach for leveraging the wealth of social media documents available on the Web for event identification and characterization. To match event occurrences in social media, we develop a new visual-based method for retrieving events in huge photo collections, typically in the context of User Generated Content. The main contributions of the thesis are the following : (1) a new visual-based method for retrieving events in photo collections, (2) a scalable and distributed framework for Nearest Neighbors Graph construction for high dimensional data, (3) a collaborative content-based filtering technique for selecting relevant social media documents for a given event.

KEY-WORDS : Event matching, event mining, event summarization, nearest neighbors graph.

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